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TITLE OF INVENTION NETWORK FORWARDING DEVICE AND NETWORK NEXT-HOP SELECTION METHOD

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Applicant herewith submits to the United States Designated/Elected Office (DO/EO/US) the following items and other information:

1. ☒ This is a **FIRST** submission of items concerning a filing under 35 U.S.C. 371.
2. ☐ This is a **SECOND** or **SUBSEQUENT** submission of items concerning a filing under 35 U.S.C. 371.
3. ☐ This express request to begin national examination procedures (35 U.S.C. 371(f)) at any time rather than delay examination until the expiration of the applicable time limit set in 35 U.S.C. 371(b) and PCT Articles 22 and 39(1).
4. ☒ A proper Demand for International Preliminary Examination was made by the 19th month from the earliest claimed priority date.
5. ☒ A copy of the International Application as filed (35 U.S.C. 371(c)(2))
 - a. ☐ is transmitted herewith (required only if not transmitted by the International Bureau).
 - b. ☒ has been transmitted by the International Bureau.
 - c. ☐ is not required, as the application was filed in the United States Receiving Office (RO/US).
6. ☒ A translation of the International Application into English (35 U.S.C. 371(c)(2)).
7. ☐ Amendments to the claims of the International Application under PCT Article 19 (35 U.S.C. 371(c)(3))
 - a. ☐ are transmitted herewith (required only if not transmitted by the International Bureau).
 - b. ☐ have been transmitted by the International Bureau.
 - c. ☐ have not been made; however, the time limit for making such amendments has NOT expired.
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8. ☐ A translation of the amendments to the claims under PCT Article 19 (35 U.S.C. 371(c)(3)).
9. ☐ An oath or declaration of the inventor(s) (35 U.S.C. 371(c)(4)).
10. ☐ A translation of the annexes to the International Preliminary Examination Report under PCT Article 36 (35 U.S.C. 371(c)(5)).

Items 11. to 16. below concern document(s) or information included:

11. ☐ An Information Disclosure Statement under 37 CFR 1.97 and 1.98.
12. ☐ An assignment document for recording. A separate cover sheet in compliance with 37 CFR 3.28 and 3.31 is included.
13. ☐ A **FIRST** preliminary amendment.
☐ A **SECOND** or **SUBSEQUENT** preliminary amendment.
14. ☐ A substitute specification.
15. ☒ A change of power of attorney and/or address letter.
16. ☒ Other items or information:

PCT Request Form

International Search Report with References

Information Disclosure Sheet Under 37 CFR 1.56 with References

Figs. 1-39

Credit Card Payment Form

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PCT/JP98/01232

17. ☒ The following fees are submitted:**BASIC NATIONAL FEE (37 CFR 1.492 (a) (1) - (5)) :**

Neither international preliminary examination fee (37 CFR 1.482)
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CLAIMS	NUMBER FILED	NUMBER EXTRA	RATE		
Total claims	10 - 20 =	0	X \$18.00	\$	0.00
Independent claims	4 - 3 =	1	X \$78.00	\$	78.00
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NOTE: Where an appropriate time limit under 37 CFR 1.494 or 1.495 has not been met, a petition to revive (37 CFR 1.137(a) or (b)) must be filed and granted to restore the application to pending status.

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DESCRIPTION

NETWORK FORWARDING DEVICE AND NETWORK

NEXT-HOP SEARCH METHOD

TECHNICAL FIELD

The present invention relates to a network forwarding device such as a router on a computer network system. More particularly, the present invention relates to a network forwarding device and a network next-hop search method for quickly searching for an address to which a packet, received by a network forwarding device, is to be forwarded based on the destination address contained in the packet.

10 BACKGROUND ART

In a network system, network forwarding devices such as routers are used to connect a plurality of subnets. A router checks the destination address of a packet received from a subnet to which the router is connected, determines the next hop, and forwards the received packet to the subnet to which the next-hop router or the host is connected. FIG. 1 shows the configuration of a general network system composed of a plurality of subnets connected by routers.

20 In FIG. 1, R1 and R2 are routers, SN1 is a sub network connected to port P11 of router R1, SN2 is a sub network connected to port P12 of router R1 and port 21 of router R2, H10 and H11 are hosts connected

to subnet SN1, H20 and H21 are hosts connected to subnet SN2, and H30 and H31 are hosts connected to subnet SN3.

When a packet is sent from host H10 to host
5 H21, router R1 checks the destination address DA stored in the packet as the header information to find that destination host H21 is on subnet SN2 and that subnet SN2 is connected directly to router R1. Then, router R1 outputs the packet to port P12 which connects to
10 subnet SN2. When outputting the packet, router R1 sets the destination host (H21) as the address to which the packet to is be forwarded next (next hop address).

When a packet is sent from host H10 to host
H31, router R1 checks the destination address DA stored
15 in the packet as the header information to find that destination host H31 is on subnet SN3 and that subnet SN3 is connected not directly to router R1 but via router R2. Router R1 outputs the packet to port P21 connected to subnet SN2 to which router R2 is
20 connected. When outputting the packet, router R1 sets router R2 as the address to which the packet to is be forwarded next (next hop address). Upon receiving the packet, router R2, like router R1, checks the destination address DA and forwards the packet to host
25 H31.

Next, referring to FIG. 2, the following describes the search specification used by a router, which has received a packet, to search for the address

to which the packet is to be forwarded next and the port to which the packet is to be output. TBL, a path search table, contains configuration definition information entered manually and information acquired
5 through communication of connection information between routers.

With the set of a sub-network address and a sub-network mask length as the search key, the path search table TBL is used to search for an output port,
10 next hop address, and information indicating whether or not the sub-network is connected directly (hereafter called next hop information).

In the path search specification, the mask beginning with the most significant bit and extending
15 for the number of bits indicated by the sub-network mask length is multiplied by the destination address for comparison with a sub-network address. As a result of comparison, multiple entries E1, E2, and E4 with different mask lengths match. Out of the matching
20 entries, the next hop information (next hop 2) of E2 with the longest mask is selected as the search result.

The Radish algorithm is known as one of the quick search methods in accordance with this search specification. The Radish algorithm is described, for
25 example, in "Let's read kernel (8) Path control mechanism at IP layer (2)" by Hide Yamaguchi, pp. 20-25 in UNIX MAGAZINE, April, 1997.

The following describes the Radish algorithm.

The Radish algorithm maps the nodes of a tree structure composed of multiple pointer-connected nodes, each having the left and right pointers, to the path entries. When this tree is followed, the left or right
 5 pointer of each node is selected to move to the next node until the node to which the desired path entry is mapped is reached.

First, referring to FIG. 3, the tree structure will be described. Because the concept of
 10 the tree structure does not depend on the number of bits, the address is three bits in length in FIG. 3 for easy understanding.

As shown in FIG. 3, the nodes are called, beginning with the top of the tree, 0-bit mask node, 1-
 15 bit mask nodes, 2-bit mask nodes, and 3-bit mask nodes.

In the 0-bit mask node N0000, one of the left and right pointers is selected according to whether bit 0 of the destination address is 0 or 1 to move down to 1-bit mask node N0001 or N1001. In a 1-bit mask node,
 20 one of the left and right pointers is selected according to whether bit 1 is 0 or 1 to move down to 2-bit mask node N0002, N0102, N1002, or N1102. In a 2-bit mask node, one of the left and right pointers is selected according to whether bit 2 is 0 or 1 to move
 25 down to 3-bit mask node N0003, N0013, N0103, N0113, N1003, N1013, N1103, or N1113.

When searching for a desired destination address through pointers, beginning at 0-bit mask node

N0000 of the tree, according to whether the bit is 0 or 1, the 0-bit mask node is always selected. 1-bit mask node N0001 or N1001 is selected, beginning with the left end, when the bits of the destination address is 0XX or 1XX, respectively. 2-bit mask node N0002, N0102, N1002, or N1102 is selected, beginning with the left end, when the bits of the destination address is 00X, 01X, 10X, or 11X, respectively. 3-bit mask node N0003, N0013, N0103, N0113, N1003, N1013, N1103, or N1113 is selected, beginning with the left end, when the bits of the destination address is 000, 001, 010, 011, 100, 101, 110, or 111, respectively. X indicates a 'don't care' bit, that is, any bit value, 0 or 1, is acceptable.

Therefore, 0-bit mask node N0000 is selected when the destination address belongs to the sub-network address 000/0. 1-bit mask node N0001 or N1001 is selected when the destination address belongs to the sub-network address 000/1 or 100/1. 2-bit mask node N0002, N0102, N1002, or N1102 is selected when the destination address belongs to the sub-network address 000/2, 010/2, 100/2, or 110/2. 3-bit mask node N0003, N0013, N0103, N0113, N1003, N1013, N1103, or N1113 is selected when the destination address belongs to the sub-network address 000/3, 001/3, ..., or 111/3. The notation "sss/m" indicates that "sss" represents a sub-network address and m represents the mask length.

As described above, the nodes of this tree

correspond, one to one, with all subnets each with a unique sub-network address and a mask length.

Nodes N0000, N0013, N0102, N1001, and N1103, which correspond to the entries of the path table shown in FIG. 4, are attached "*" in the tree. When the destination address DA011 is searched for in this tree by checking the pointers if each bit is 0 or 1, it is found that nodes N0000 and N0102, which have "*" attached, correspond to the matching entries. When there are a plurality of path table entries that match the destination address, the rule is that the longest mask-length sub-network is selected. That is, out of matching nodes N0000 and N0102 to which "*" is attached, the path information allocated to node N0102 nearest to the end point is selected as the search result of the path table.

As shown in the search method described above, those nodes to which "*" is not attached and which are not intermediate paths to a node with "*" attached, that is, N0003, N0103, N0113, N1003, N1013, N1113, and N1002, do not affect the search result if removed from the tree. Rather, when "*" is not attached to the bottom-level node, the search becomes more efficient because the search ends before reaching the bottom level. FIG. 5 shows a tree created by removing nodes to which "*" is not attached and which are not intermediate path to a node to which "*" is attached.

FIG. 7 shows a 2-branch (binary) tree which corresponds to the 32-bit address path table shown in FIG. 6 and which is created by the method described above. As shown in the figure, a long string NSI of 5 nodes with neither a branch nor "*" is created. The following describes the quick search method in which only one of left and right pointers of a node connects to the next node and the nodes to which path entries are not mapped are removed.

10 In this quick search method, node string NS1 with no branch or "*" is removed and the immediate higher-level node N000000000 of the removed node string NS1 is connected to the immediate lower-level node N8504000015 in the branch direction (right side in FIG. 15 7). The resulting tree is shown in FIG. 8. Removing the intermediate nodes like this is called the reduction of a tree.

Next, the path search method for a reduced tree will be described.

20 In the example shown in FIG. 8, immediately after bit 0 of the 0-bit mask node N000000000 is checked, 15-bit mask node N8504000015 is checked. Therefore, even if bit 15 of 15-bit mask node N8504000015 is checked, the intermediate bits, bit 1 to 25 bit 14, cannot be checked. To check bits 1 to 15 in one search operation, bits 1 to 15 of the destination address are compared with bits 1 to 15 of the sub-network address 0x85.04.00.00 of node N8504000015. An

equal comparison indicates that the intended node is reached, that is, the node would be reached also by comparing them, one bit at a time, in a non-reduced tree. An unequal comparison indicates that a non-
5 intended node is reached, that is, there would be no destination in the non-reduced tree.

In FIG. 8, an equal condition always occurs in bit 0 because the branch is already selected based on the equality. In general, if a check is made each
10 time a node is reached to see if the correct node is reached, it is guaranteed that the destination address equals the sub-network address of the node for the mask length beginning with bit 0. Thus, when the next node is reached, whether the correct node is reached is
15 checked by comparing the destination address with the sub-network address for the mask length beginning with bit 0 regardless of how many bits were tested before.

DISCLOSURE OF THE INVENTION

As described above, the problem with the
20 Radish algorithm is that it takes long because the algorithm checks destination addresses from the top downward, one bit at a time, to search for a path.

It is an object of the present invention to provide a network forwarding device, especially a
25 router, that executes quick path search processing when transferring packets.

It is another object of the present invention

to provide a network next-hop search method, for use in a network forwarding device such as a router, that checks the destination address stored in a received packet to quickly search for a packet forwarding
5 address.

To achieve the above object, the method according to the present invention expands several high-order bits of a sub-network address into fixed positions in the memory during the search for a part
10 corresponding to the high-order bits of the sub-network address. This method eliminates the time required for the search for the several high-order bits of the sub-network address.

In addition, several high-order bits of the
15 nodes of a search tree are contained in a search processing LSI. Search processing in the internal memory of the LSI and search processing in the external memory are performed in pipeline mode to hide the time required for searching for several high-order bits.

20 In addition, the number of branches of a node of a search tree, which is two branches in the prior art node, is increased in increments of a power of 2, such as 4-branch tree nodes, 8-branch tree nodes and more. This method increases the number of bits to be
25 tested at a time from one bit for each node to continuous two, three, and more bits, thereby decreasing the number of nodes to be followed until the end of search.

In addition, to reduce the amount of memory for storing a data structure representing a search tree when configuring a 4-branch tree, an 8-branch tree, and more in general, a 2^p -branch tree, a total of (2^p-1) 2-branch tree nodes, composed of one 2-branch tree node and immediately lower 2-branch tree nodes of $(p-1)$ levels, are combined into one 2^p -tree node and, into the combined lowest-level $2^{(p-1)}$ 2-branch tree nodes, path data allocated to higher-level nodes is embedded to form the 2^p -tree node with $2^{(p-1)}$ 2-branch tree nodes. Furthermore, only one copy of a sharable element is provided.

In addition, when a 4-branch tree node, 8-branch tree node, or more-branch tree node composed of a plurality of 2-branch tree nodes is read for search, only a required part, not the whole, of the node is read to prevent the read time from increasing as the node becomes larger. To select only the required part of node data, node mask length information is necessary before the node data is read. Thus, each node contains the mask length of the immediately-lower node. Furthermore, at the start of each node, a flag indicating whether or not a path is allocated to the node is provided. This flag is read first to prevent path information for a node, to which a path is not allocated, from being read. This reduces the data read time.

BRIEF DESCRIPTION OF THE DRAWINGS

Fig. 1 is a diagram showing the configuration of a general network system to which the present invention applies. FIG. 2 is a diagram showing a path search specification for a router. FIG. 3 is a diagram showing a 2-branch tree which has all nodes when the address is three bits long. FIG. 4 is a diagram showing an example of a path table when the address is three bits long. FIG. 5 is a diagram showing a tree from which a node which has no path and which is not an intermediate path to another node with a path is removed. FIG. 6 is a diagram showing an example of a path table when the address is 32 bits long. FIG. 7 is a diagram showing a tree corresponding to the path table shown in FIG. 6. FIG. 8 is a diagram showing a tree from which a node which has no branch and to which no path is allocated is removed. FIG. 9 is a diagram showing a tree whose 2-bit mask node is expanded in memory to eliminate the search time for bits 0-1. FIG. 10 is a diagram showing the configuration of memory in which nodes with a k-bit or shorter mask are stored in a path search LSI. FIG. 11 is a diagram showing the timing chart of conventional path search processing where pipeline processing is not performed. FIG. 12 is a diagram showing the timing chart of the pipeline processing of path search when the nodes with a k-bit or shorter mask are stored in a path search LSI. FIGS. 13-15 are diagrams showing 2-branch tree nodes to be

transformed into one 8-branch tree node that are enclosed by dotted lines. FIGS. 16 and 17 are diagrams showing the trees each of which is used to search for bits to the bit division position by expanding the first node into memory when bit positions are divided such that the first division does not start with mask length 0. FIGS. 18-20 are diagrams showing trees that eliminate the search time of more starting bits by combining the method shown in FIGS. 16-17 and the method shown in FIG. 9. FIG. 21 is a diagram showing a tree created by transforming three 2-branch tree nodes, which are combined into one 4-branch tree node, into two 2-branch tree nodes when transforming a 2-branch tree to a 4-branch tree. FIGS. 22-30 are diagrams showing how three 2-branch nodes, which are combined into one 4-branch tree node, are transformed into two 2-branch tree nodes. FIG. 31 is a diagram showing how seven 2-branch nodes, which are combined into one 8-branch tree node, are transformed into four 2-branch tree nodes. FIG. 32 is a diagram showing a model of tree expansion when the memory amount of a tree is estimated. FIG. 33 is a diagram showing the comparison between the memory amount required when the tree is configured with 2-branch tree nodes and the memory amount required when the tree is configured with 4- and 8- branch tree nodes, with the node presence probability taken into consideration. FIG. 34 is a diagram showing the structure of a 2-branch tree node.

FIG. 35 is a diagram showing the structure of a 4-branch tree node. FIG. 36 is a diagram showing a method for reading only a part, not the whole, of a node to prevent an increase in the node data read time when the node becomes large. FIG. 37 is a diagram showing a method for not reading an element that is not necessary depending upon a condition to reduce the read time and to increase processing speed. FIG. 38 is a block configuration diagram showing an example of the configuration of a router. FIG. 39 is a block diagram showing an another example of the configuration of the router.

BEST MODE FOR CARRYING OUT THE INVENTION

To describe the present invention more in detail, the best mode for carrying out the present invention will be described with reference to the attached drawings.

First, the typical configuration of a router to which the present invention applies will be described with reference to FIG. 38. In FIG. 38, number 100 is a router, number 110 is a routing controller, number 120 is a router bus, number 130 is a network interface, number 140 is a port, and number 150 is a sub-network.

The network interface 130 receives a packet from a sub-network connected to the port 140 and sends the received packet to the routing controller 110 via

the router bus 120. The routing controller 110, with a routing table in which routing information is stored, checks the destination of the received packet with the use of the routing information to determine the sub-

5 network 150 to which the packet is to be forwarded, and sends the packet to the network interface 130 of the port 140 to which the sub-network 150 is connected. Upon receiving the packet from the routing controller 110, the network interface 130 sends the packet to the

10 sub-network 150 to which the packet is to be forwarded. The routing controller 110 updates and maintains the routing information stored in the routing table based on the header information of the received packet and, at the same time, manages the whole router 100.

15 FIG. 39 is a block diagram showing another example of the configuration of the router. In FIG. 39, number 200 is a router, number 210 is a routing processor (RP), number 220 is intra-router communication means, number 230 is a network interface,

20 number 240 is a port, number 250 is a sub-network, and number 260 is a router manager. In this configuration, the routing controller 110 shown in FIG. 38 is divided into two: the routing processor 210 that executes the routing function and the router manager 260 that

25 manages the router 200. At the same time, there are a plurality of components each composed of one or more network interfaces 230, which are shown also in FIG. 38, and the routing processor 210. The router manager

260 manages the whole router 200 and, at the same time, distributes routing information to the routing processors 210. The intra-router communication means 220, a crossbar switch or a bus, performs communication among the routing processors 210 or between the routing processor 210 and the router manager 260. The routing processor 210, like the routing controller 110 shown in FIG. 38, transfers packets with the network interface 230 connected thereto. In addition, when the routing processor 210 transfers a packet to the sub-network 250 connected to some other routing processor 210, it transfers the packet to the corresponding routing processor 210 via the intra-router communication means 220.

Next, next-hop path search processing executed by the routing controller 110 and the routing processor 210 will be described. First, the high-speed method of the Radish algorithm according to the present invention will be described.

The first high-speed method will be described with reference to FIG. 9. When the tree is not reduced, the conventional Radish algorithm searches the tree for a path beginning at the 0-bit mask node, one bit at a time. On the other hand, the method according to the present invention expands the m-bit mask nodes in fixed positions in the memory, as shown in FIG. 9, regardless of whether m-bit mask nodes are present or not. FIG. 9 shows an example in which the 2-bit mask

nodes, N0002, N0102, N1002, and N1102, are expanded in fixed positions in the memory.

In this case, the conventional method starts with 0-bit mask node N0000, jumps to one of 1-bit mask nodes N0001 and N1001 according to the value of bit 0, and reaches one of 2-bit mask nodes N0002 and N0102 or one of N1002 and N1102 according to the value of bit 1.

The method according to the present invention finds the address where 2-bit mask nodes N0002, N0102, N1002, and N1102 are expanded using the values of bits 0 and 1 and then jumps directly to one of 2-bit mask nodes N0002, N0102, N1002, and N1102. This reduces the search time by the period of time required for two node searches.

In general, jumping directly to a fixed position in the memory where the m-bit mask nodes are expanded reduces the search time by the time required for traversing m nodes from the 0-bit mask node to (m-1)bit mask nodes. On the other hand, this method decreases memory efficiency because the 2^m m-bit mask nodes must be expanded in memory regardless whether or not there are such nodes. Therefore, the value of m should be decided considering the trade-off between memory efficiency and performance.

Next, the second high-speed method will be described with reference to FIGS. 10 to 12. As shown in FIG. 10, the nodes with a k-bit or shorter mask are expanded in memory ML1 contained LSI L1 in which path

search is executed, while the nodes with a $(K+1)$ bit or longer mask are expanded in memory M1 external to LSI L1. Because access to memory M1 in LSI L1 is fast and because internal memory ML1 of LSI L1 is independent of
 5 external memory M1, pipeline processing is possible.

This pipeline processing speeds up address search processing. In general, although the memory amount of LSI L1 is smaller than that of external memory, there is a small number of shorter-mask nodes. Therefore,
 10 the shorter-mask nodes may be contained in LSI L1.

FIG. 11 shows the timing chart of the conventional method in which no pipeline processing is executed. FIG. 12 shows the timing chart of path
 15 search processing according to the present invention that is executed in pipeline mode between the internal memory ML1 of LSI L1 and external memory M1.

Conventionally, node search processing for packets (packets 1, 2, 3) is executed sequentially, as shown in FIG. 11. That is, node search processing PR10 for a k -
 20 bit or shorter mask and node search processing PR11 for a $(k+1)$ bit or longer mask for packet 1, PR20 and PR21 for packet 2, and PR30 and PR31 for packet 3 are executed sequentially. On the other hand, the method according to the present invention is executed as shown
 25 in FIG. 12. That is, node search processing PR10 for a k -bit or shorter mask is executed for packet 1 using the internal memory ML1 of LSI L1 and, then, node search processing PR11 for a $(k+1)$ bit or longer mask is

started for packet 1 using the external memory M1 of LSI L1. At the same time, node search processing PR20 for a k-bit or shorter mask is executed for packet 2 using the internal memory ML1 of LSI L1. After that, 5 when processing PR20 and PR11 are finished, node search processing PR21 for a (k+1)bit or longer mask is executed for packet 2 using the external memory M1 of LSI L1. At the same time, node search processing PR20 for a k-bit or shorter mask is executed for packet 3 10 using the internal memory ML1 of LSI L1. Thereafter, packet search processing is executed similarly in pipeline mode.

The third high-speed method provides 2^p branches for one node to search for p bits at a time, 15 in contrast to two branches for one node in the conventional method to search for one bit at a time. Thus, this method reduces the conventional search time by $1/p$. In the description below, a node with 2^p branches is called a 2^p -branch tree node.

20 A 2^p -branch tree node is created by transforming a tree composed of conventional binary (2-branch) tree nodes. To transform a tree, one n-th bit 2-branch tree node and the (n+1)th to (n+p-1)th bit 2-branch tree nodes below the 2-branch tree node are made 25 to correspond to a 2^p -branch tree node. As an example, the transformation from a 2-branch tree to an 8-branch tree is shown in FIGS. 13 to 15.

For an 8-branch tree, the bit positions of 2-

branch tree nodes which are made to correspond to an 8-branch tree node is one of the following three:

(a) As shown in FIG. 13, 0-bit to 2-bit nodes, 3-bit to 5-bit nodes, 6-bit to 8-bit nodes, 9-bit to 11-bit nodes, 12-bit to 14-bit nodes, 15-bit to 17-bit nodes, 18-bit to 20-bit nodes, 21-bit to 23-bit nodes, 24-bit to 26-bit nodes, 27-bit to 29-bit nodes, and 30-bit to 32-bit nodes are each made to correspond to one 8-branch tree node.

(b) As shown in FIG. 14, 1-bit to 3-bit nodes, 4-bit to 6-bit nodes, 7-bit to 9-bit nodes, 10-bit to 12-bit nodes, 13-bit to 15-bit nodes, 16-bit to 18-bit nodes, 19-bit to 21-bit nodes, 22-bit to 24-bit nodes, 25-bit to 27-bit nodes, 28-bit to 30-bit nodes, and 31-bit to 32-bit nodes are each made to correspond to one 8-branch tree node.

(c) As shown in FIG. 15, 2-bit to 4-bit nodes, 5-bit to 7-bit nodes, 8-bit to 10-bit nodes, 11-bit to 13-bit nodes, 14-bit to 16-bit nodes, 17-bit to 19-bit nodes, 20-bit to 22-bit nodes, 23-bit to 25-bit nodes, 26-bit to 28-bit nodes, 29-bit to 31-bit nodes, and 32-bit node are each made to correspond to one 8-branch tree node.

The tree may be configured in any bit division. To make path addition and deletion easy, one of the above three bit position divisions is used.

Out of the above three bit position divisions, the mask does not begin with bit 0 in the

divisions other than the first. This means that the search for the first bit must be made separately. In this search, the method shown in FIG. 9, in which m-bit mask nodes are expanded into memory, is used. FIGS. 16 and 17 show the configuration of the trees whose first bits of the 8-branch tree nodes corresponding to the division positions corresponding to FIGS. 14 and 15 are expanded into memory.

In the configuration shown in FIG. 16, 1- to 3-bit mask nodes N8013 and N8113 are arranged in fixed positions in memory. One of them is selected according to whether bit 0 contains 0 or 1. In the configuration shown in FIG. 17, 2- to 4-bit mask nodes N80024, N80124, N81024, and N81124 are arranged in fixed positions in memory. One of them is selected according to whether bits 0-1 contain 00, 01, 10, or 11.

Or, the method shown in FIG. 9 in which the m-bit masks are expanded into memory may be combined with the method in which 2^p -branch tree nodes are created to search for multiple number of bits at a time. More specifically, instead of arranging 1, 2, and 4 nodes from top to bottom in the bit division positions shown in FIGS. 13, 14, and 15, 8 times the number of nodes (8, 16, and 32), 8^2 times the number of nodes (64, 128, 256), or, more in general, 8^q times the number of nodes may be used to eliminate the need for the first, second, or q-th search time of 2^p branch nodes. When $p=3$ and $q=1$, that is, when one node-search

time is eliminated in the 8-branch tree nodes, FIGS. 18, 19, and 20 show node expansion in memory for the three bit division positions shown in FIGS. 13, 14, and 15. Increasing the values of p and q, like this, 5 reduces the path search time but requires more memory. The values of p and q should be decided based on the trade-off between memory efficiency and performance.

Next, the method for configuring 4-branch tree nodes, 8-branch tree nodes, 16-branch tree nodes, 10 or more in general, power-of-2-branch tree nodes will be described with reference to FIG. 21.

FIG. 21 shows a 4-branch tree. In this example, the 2-branch nodes, each composed of a total of three 2-branch tree nodes, that is, one 2-branch 15 tree nodes A, B, C, D, or E and each respective two immediately lower 2-branch tree nodes A0, A1, B0, B1, C0, C1, D0, D1, E0, E1, are combined into one of 4-branch tree nodes N401, N4023, N4123, N4223, and N4323. A total of three 2-branch tree nodes are transformed 20 into a node with the size of lower-level 2-branch tree nodes. This transformation is performed according to the path search specification, that is, when a match occurs on a plurality of paths each having its own mask length, the path with the longest mask length is used. 25 In addition, this transformation is performed so that, when path search is made, the 2-branch path search matches the 4-branch path search result.

The node transformation of 4-branch tree

nodes is shown in FIGS. 22-30.

When all three nodes are present (FIG. 22) and path information is allocated to all nodes, the path information *A0 and *A1 of the lower nodes are not
 5 deleted but the path information *A of the higher node is deleted. This is because, when a match occurs on the path of node A, a match always occurs on the path of node A0 or node A1. Therefore, *A is never used according to the path search specification which
 10 stipulates that, when a match occurs on a plurality of paths each having its own length, the path with the longer mask length is used.

When path information *A is allocated to higher node A but path information is not allocated to
 15 A1 of lower nodes A0 and A1 (FIG. 23), the path information *A of A is given to A1 as the path information. This also applies when a path is not allocated only to A0 of two lower nodes A0 and A1.

When a path is not allocated to both lower
 20 nodes A0 and A1 (FIG. 24), the path information A* of A is given to both A0 and A1 as the path information.

When there is no lower node A1 (FIG. 25), the corresponding node is supplied and the path information *A of higher node A is given to the supplied node.
 25 Because no lower node is connected node A1, NULL is set in the pointer to the lower node of node A1. This also applies when a path is not allocated only to A0 of two lower nodes A0 and A1.

When there is no lower node, neither A0 nor A1 (FIG. 26), both nodes are supplied, the path information *A of A is given to them, and NULL is given to the pointers to the nodes below them.

5 When a path is not allocated to the higher node (FIG. 27), the higher node is discarded.

 When path information *A is not allocated to node A and, for the lower nodes A0 and A1, a path is not allocated to A1 (FIG. 28), A1 has no path
10 information even in a 4-branch node tree. This also applies when a path is not allocated to A0 of lower nodes A0 and A1.

 When a path is not allocated to both lower nodes A0 and A1 (FIG. 29), there is no path information
15 for both nodes in the 4-branch tree node.

 When there is lower node A0 only (FIG. 30), lower node A1 is supplied. This also applies when there is lower node A1 only.

 The same transformation is performed also for
20 an 8-branch tree. That is, a total of seven 2-branch tree nodes are transformed into a node with the size of four lowest-level 2-branch tree nodes. FIG. 31 shows two examples of transforming three higher 2-branch tree nodes.

25 FIG. 31(a) shows an example in which all seven 2-branch tree nodes to be combined into one are present but path information is not allocated to some of them. Out of four lowest nodes, path information is

not allocated to A01 and A10. To these nodes, the path information *A and *A1, which are their nearest higher nodes to which path information is allocated, that is, the path information with the longest mask nodes (A and A1, respectively), is given.

FIG. 31(b) shows an example in which some of seven 2-branch tree nodes to be combined into one are not present. In this case, missing nodes A01 and A10 are first supplied as nodes to which path information is not allocated and then path information is given to them according to the rule used in the example in FIG. 31(a). Because, out of lowest four nodes A00, A01, A10, and A11, the supplied nodes A01 and A10 have no lower nodes, NULL is set in the pointers to the lower nodes.

In general, the method described above is also used for a 2^p -branch tree. (2^p-1) 2-branch tree nodes are transformed into a node with the size of the lowest-level 2^{p-1} nodes.

Reducing the size of a node in this manner prevents memory efficiency from decreasing. In the following example, the memory amount required when a tree is created as a 2-branch tree under the approximation condition given below and the memory amount required when a tree is created as a 2^p -branch tree are calculated. This example shows that a 2^p -branch does not decrease memory efficiency, especially when p is small.

Approximation:

1. In a 2-branch tree, paths are allocated only to the end nodes.
2. All 32 bits have branches.
- 5 3. An increase in the number of branches is fixed for all bits.
4. The number of paths is about 10k.

FIG. 32 shows the form of a tree configured under the approximation condition. According to this
 10 approximation, the tree becomes $10K^{(1/32)}$ times larger, that is, about 1.33 times larger, each time the tree goes down one-bit hierarchical level.

Referring to FIG. 33(a), the memory amount required when the tree is composed of 2-branch tree
 15 nodes and the memory amount required when the tree is composed of 4-branch tree nodes are compared. As shown in the figure, N2, one of three 2-branch tree nodes constituting the 4-branch tree node, has about 1.33 immediately lower nodes. That is, the node has right
 20 and left nodes, N20 and N21, each with a probability of 0.67. Therefore, with the node presence probability taken into consideration, the average of the total amount of memory required for 2-branch tree nodes N2, N20, and N21 constituting the 4-branch tree node is the
 25 memory amount required for $(1 + 1.33)$ 2-branch tree nodes, that is, 2.33 2-branch tree nodes. When constituting the tree with 4-branch tree nodes, these three 2-branch tree nodes are combined into one 4-

branch tree node N4. The memory amount required for this one 4-branch tree node equals the memory amount required for two 2-branch tree nodes, N20 and N21.

Similarly, referring to FIG. 33(b), the
 5 memory amount required when the tree is composed of 2-branch tree nodes and the memory amount required when the tree is composed of 8-branch tree nodes are compared. As shown in the figure, N2, the highest node of seven 2-branch tree nodes N2, N20, N21, N200, N201,
 10 N210, and N211 constituting the 8-branch tree nodes, has about 1.33 immediately lower nodes, N20 and N21, which in turn have about 1.33^2 , that is, about 1.78, lower nodes, N200, N201, N210, and N211. Therefore, with the node presence probability taken into
 15 consideration, the average of the total amount of memory required for 2-branch tree nodes N2, N20, N21, N200, N201, N210, and N211 constituting the 8-branch tree nodes is the memory amount required for $(1 + 1.33 + 1.78)$ 2-branch tree nodes, that is, 4.11 2-branch
 20 tree nodes.

When constituting the tree with 8-branch tree nodes, these seven 2-branch tree nodes are combined into one 8-branch tree node N8. The memory amount required for this one 8-branch tree node equals the
 25 memory amount required for four 2-branch tree nodes, N200, N201, N210 and N211.

The following shows the memory amount comparison between a 2-branch tree and 4-, 8-, 16-, 32-,

64-, 128-, and 256-branch trees that is made in the manner described above.

2-branch tree : 4-branch tree = 2.33 : 2 = 0.86 times

2-branch tree : 8-branch tree = 4.11 : 4 = 0.97 times

5 2-branch tree : 16-branch tree = 6.48 : 8 = 1.23 times

2-branch tree : 32-branch tree = 9.65 : 16 = 1.66 times

2-branch tree : 64-branch tree = 13.86 : 32 = 2.31 times

2-branch tree : 128-branch tree = 19.49 : 64 = 3.28 times

2-branch tree : 256-branch tree = 26.98 : 128 = 4.74 times

10 When the number of paths is 1M, the tree becomes $1M^{(1/32)}$ times larger, that is, about 1.54 times larger, each time the tree goes down one-bit hierarchical level. The memory amount ratios between the 2-branch tree and the 4-, 8-, 16-, 32-, 64-, 128-,
15 and 256-branch tree are as follows:

2-branch tree : 4-branch tree = 2.54 : 2 = 0.79 times

2-branch tree : 8-branch tree = 4.91 : 4 = 0.81 times

2-branch tree : 16-branch tree = 8.56 : 8 = 0.93 times

2-branch tree : 32-branch tree = 14.19 : 16 = 1.13 times

20 2-branch tree : 64-branch tree = 22.85 : 32 = 1.40 times

2-branch tree : 128-branch tree = 36.18 : 64 = 1.77 times

2-branch tree : 256-branch tree = 56.72 : 128 = 2.26 times

 The result is that, under the above assumption, memory usage efficiency is increased for up
25 to an 8-branch tree when the number of paths is 10k, and for up to a 16-branch tree when the number of paths is 1M. Even for a 256-branch tree, the required memory amount is 3.28 times that of a 2-branch tree when the

number of paths is 10k, and 2.26 times when the number of paths is 1M. The memory usage efficiency remains high for the following reasons:

(1) Combining p-level nodes into one makes the
5 memory of combined nodes more compact than the memory of nodes before being combined. That is:

Three 2-branch tree nodes become a 4-branch tree node with the size two times the 2-branch tree node.

Seven 2-branch tree nodes become an 8-branch tree node
10 with the size four times the 2-branch tree node.

15 2-branch tree nodes become a 16-branch tree node with the size eight times the 2-branch tree node.

31 2-branch tree nodes become a 32-branch tree node with the size 16 times the 2-branch tree node.

15 63 2-branch tree nodes become a 64-branch tree node with the size 32 times the 2-branch tree node.

127 2-branch tree nodes become a 128-branch tree node with the size 64 times the 2-branch tree node.

255 2-branch tree nodes become a 256-branch tree node
20 with the size 128 times the 2-branch tree node.

(2) The tree expands as the tree goes down one bit level, increasing the usage efficiency of data in one node (As the number of supported paths increases, the tree expands more and usage efficiency of data in
25 one node increases).

In addition, a 4-branch tree node, 8-branch tree node, 16-branch tree node, and so on, each handle 2, 4, 8, ... 2-branch tree nodes at a time. Thus, for

an element that may be shared among the combined 2-branch tree nodes, only one such element is required. This reduces the amount of memory required for a 4-branch tree node, 8-branch tree node, 16-branch tree node, and so on. Examples of an element that may be shared among combined 2-branch tree nodes are a sub-network address and a sub-network mask length. As will be described later, the sub-network mask length is not the sub-network mask length of the node itself but that of the node immediately below the node. Therefore, the memory amount is not reduced much.

FIGS. 34 and 35 show, respectively, the structure of a 2-branch tree node and the structure of a 4-branch tree node which is generated by combining two 2-branch tree nodes and which has only one network address that may be shared between two combined 2-branch tree nodes. FIG. 34 is a diagram showing the structure of a 2-branch tree node. The next node mask length 0 field and the next node mask length 1 field contain, not the sub-network mask length of this node, but the sub-network mask length of the nodes immediately below this node. The mask lengths of the nodes immediately below, not that of the node itself, are set in those fields for fast access. This will be described later with reference to FIG. 36. Flag 0 and Flag 1 each contain a bit indicating whether the sub-network corresponding to this node connects directly to the router or via one or more routers. Also, as will

be described later with reference to FIG. 37, those flags contain information indicating, for example, whether or not a path is allocated to this node, that is, whether or not "*" is attached to this node as in the example of the tree shown in FIG. 8. Flag 0 and Flag 1 contain the same value. This is because only one of word W0 and word W1 need be read. Speeding up processing by reading only a part of a node, instead of reading the entire node, will be described later with reference of FIG. 36. The pointer 0 to next node field and the pointer 1 to next node field each contain a pointer to the next node to be reached when the value of the bit position of the destination address indicated by the mask length of this node is 0 or 1.

15 The sub-network address field contains the sub-network address corresponding to the node. The output port number field and the next-hop address field contain path information allocated to this node. The next-hop address field contains the address of the port from which the received packet is to be output, and output port number field contains the address of the router to which the packet on that port is to be sent.

A 4-branch tree node is structured to have two 2-branch tree nodes and to have only one copy of data that may be shared. The sharable data is the sub-network address only. FIG. 35 shows the structure of a 4-branch tree node created in this method.

In the examples shown in FIGS. 34 and 35, a

2-branch tree node is a little bit larger than 16 bytes that is a power of 2. A 4-branch tree node, with only one sub-network address in one node, allows the node to be contained in 32 bytes that is a power of 2. An 8-
5 branch tree node, with only one sub-network address in one node, allows the node to be contained in 64 bytes and makes an extra 4-byte field available for use in storing other information. Increasing the number of 2-branch tree nodes that are combined into one further
10 reduces the size of a node with respect to the magnitude of a power of 2.

As described above, when the size of one node is within a power of 2, the hardware configuration becomes very simple. Some examples in which the
15 hardware configuration becomes simple, as well as their benefits, will be described.

As the first example, consider a case in which a 4-branch tree node is contained in 32 bytes. In this case, even when the memory is composed of a
20 plurality of banks, a one-node memory area does not extend across the bank boundary. In addition, even when a dynamic RAM is used as the memory, a one-node memory area does not extend across the row address boundary.

25 As the second example, when finding the address of an element in a node, the value of the pointer to the node need not be added to the offset from the address pointed to by the pointer. Instead,

the value in the high-order bits of the address are used as the pointer to the node, and the value of the low-order bits are used as the offset. For example, when a 4-branch tree node is contained in 32 bytes, the address of an element in a node may be represented by allocating the pointer to the node in the 2^5 or higher bit positions and allocating the address of the offset of the element within the node in the 2^0 to 2^4 bit positions.

As the third example, when a 4-branch tree node is contained, for example, in 32 bytes, the start address of the next node divided by 32 may be contained in each node as the pointer to the next node. This reduces 5 bits for each pointer in one node.

One of the problems with this method is that, when 4-branch tree nodes, 8-branch tree nodes, 16-branch tree nodes, and so on are used, the size of each node increases. The increase in the size of a node increases the time required to read one node into the search processing LSI during search time, degrading performance. This problem may be solved by reading only a part of the node instead of reading the whole node. This method will be described with reference to FIG. 36.

FIG. 36 shows an example of a 4-branch tree node. As shown in FIGS. 34 and 35, a m-bit mask 4-branch tree node has a form established by combining a 2-branch tree node whose bit m of the destination

address is 0 and a 2-branch tree node whose bit m of the destination address is 1. Therefore, this method, which checks bit m of the destination address and then reads only the corresponding 2-branch tree node, reads only the data amount equivalent to that of a 2-branch tree node even if the node size increases. In this case, as shown in Fig. 35, the sub-network address which is shared by the combined 2-branch tree nodes is read regardless of the value of bit m of the destination address.

In addition, the method checks the value of bit $m + 1$ of the destination address to read only one of two next-node pointers contained in the 2-branch tree node, further decreasing the amount of data to be read.

This method applies also to a 2-branch tree node. For example, for a m -th bit 2-branch tree node, the method checks the value of bit m of the destination address and reads only one of two next-node pointers.

After the method described above is executed and the node mask length is m , data is read in order of $(W0 \rightarrow W4 \rightarrow W5 \rightarrow W6)$, $(W1 \rightarrow W4 \rightarrow W5 \rightarrow W6)$, $(W2 \rightarrow W4 \rightarrow W5 \rightarrow W7)$, or $(W3 \rightarrow W4 \rightarrow W5 \rightarrow W7)$ according to whether the value of bits m and $m+1$ is 00, 01, 10, or 11.

As described above, reading a part of a node requires the information on the mask length m of the node. The mask length m of the node must be read first when one-node data is read, or the mask length, which

is moved to the preceding node, must be read when the preceding-node data is read. In the method in which the node mask length m is read first when one-node data is read, there are two delay factors: one is the gate
 5 delay of the search processing LSI for extracting the value of bit m of the destination address and the other is the memory read latency between the time the address of data to be read next is output to memory and the time data is read from memory into the search
 10 processing LSI. These delay factors cause a delay between the time the mask length m is read and the time data to be read next is selected and read. Therefore, reading only a part of a node in this method does not increase performance much. The method in which the
 15 node mask length m is moved to the preceding node and the preceding node is read provides better performance.

In addition, when the mask length m of a node is moved to the preceding node, one-node data should be read in order of the next-node mask length m , pointer
 20 to the next node, next sub-network address, output port number, and next hop address. This order allows the address of data to be calculated most quickly.

The pointer to the next node points to the start of the next node memory area. The offset from
 25 the start of the next node to the address of the part to be read first may be obtained by reading the mask length m of the next node and then checking the value of the corresponding bit position of the destination

address.

Next, referring to FIG. 37, the following describes how the read time is reduced and the processing speed is increased by not reading an element
 5 in one node that is not required depending upon the condition.

FIG. 37 shows an example of a 4-branch tree node. The Radish algorithm does not always allocate a path to all nodes and requires that a node be provided
 10 at a branch even if a path is not allocated thereto. As shown in FIG. 37, the flag, which is a part of node data that is read first, contains information indicating whether or not a path is allocated to the node. And, in a node to which a path is not allocated,
 15 the output port and the next hop address are not read to reduce the read time. Because the information indicating whether or not a path is allocated to the node may be represented by one bit, the increase in time required to read this information is small.

20 When the mask length of this node is m , this method reads W_0 if the value of bits m to $m + 1$ of the destination address is 00. If Flag 00 indicates that there is no path information for the 0th 2-branch constituting the 4-branch tree, the method reads only
 25 W_4 . Only when it is found that there is path information, the method reads data in order of W_4 , W_5 , and W_6 as shown in Fig. 36. This also applies when the value of bits m to $m + 1$ of the destination address is

01, 10, or 11.

INDUSTRIAL APPLICABILITY

As described above, the present invention provides a network next-hop search method for use in a
5 network forwarding device, such as a router, and a network forwarding device using the method. This method allows a network forwarding device to quickly search for the forwarding address of a received packet to increase the packet processing performance of the
10 network forwarding device.

CLAIMS

1. A network forwarding device connecting a plurality of networks comprising:

a port to which one of said networks is connected;

a network interface connected to said port and controlling an interface with the network connected to said port; and

a routing processor connected to said network interface via an intra-device communication line and performing routing of a packet received from said network interface;

wherein said routing processor includes:

path information holding means; and

next-path searching means for calculating, based on path information held in said path information holding means, a path to which the received packet is to be forwarded next and

wherein, in a 2-branch tree search in which a destination address of the received packet is checked, one bit at a time beginning at a highest-order bit, said next-path searching means combines p (p is an integer equal to or larger than 2) levels into one 2^p -branch tree to perform a search of the p levels of the 2-branch tree as a one-level search.

2. The network forwarding device according to claim 1 wherein said next-path searching means combines a total of $(2^p - 1)$ 2-branch tree nodes, composed of one

2-branch tree node and immediately lower 2-branch tree nodes of $(p-1)$ levels, into one said 2^p -tree node and, into the combined lowest-level $2^{(p-1)}$ 2-branch tree nodes, embeds path data allocated to higher-level nodes to form said 2^p -tree node with $2^{(p-1)}$ 2-branch tree nodes.

3. The network forwarding device according to claim 2 wherein, when combining a plurality of 2-branch trees, said next-path searching means has only one element if the element may be shared.

4. The network forwarding device according to claim 2 wherein said next-path searching means does not read a whole node when reading the 2^p -branch tree node but reads only data corresponding to one of the $2^{(p-1)}$ 2-branch tree nodes combined when the 2^p -tree node was created.

5. The network forwarding device according to claim 2 wherein said next-path searching means stores into each node, not a mask length of the node, but a mask length of a node immediately below the node, to find the node mask length before reading data of the node and selects a part of the node data to be read according to a value stored in a bit position, indicated by the node mask length, to the bit position + $p-1$.

6. The network forwarding device according to claim 4 wherein said next-path searching device provides a flag in data that is read first for each node, said flag indicating whether or not a path is

allocated to the node, reads the flag first, and does not read path information for a node to which the path is not allocated.

7. The network forwarding device according to claim 1 wherein said network forwarding device is a router.

8. A network forwarding device connecting a plurality of networks comprising:

a port to which one of said networks is connected;

a network interface connected to said port and controlling an interface with the network connected to said port; and

a routing processor connected to said network interface via an intra-device communication line and performing routing of a packet received from said network interface;

wherein said routing processor includes:

path information holding means; and

next-path searching means for calculating, based on path information held in said path information holding means, a path to which the received packet is to be forwarded next

wherein said next-path searching means searches for a next path using a 2-branch tree search in which a destination address is checked, one bit at a time, beginning with a highest-order bit, performs a search for a match of an address and a mask by

associating a bit position to be checked with a mask length, expands 2^m m-bit (m is a natural number) mask nodes into fixed positions in storing means, makes each of the m-bit mask nodes correspond, one to one, with a value that may be represented by bit 0 to bit (m-1) of the destination address, and selects one of the m-bit mask nodes according to a value represented by bit 0 to bit (m-1) of the destination address.

9. A network forwarding device connecting a plurality of networks comprising:

a port to which one of said networks is connected;

a network interface connected to said port and controlling an interface with the network connected to said port; and

a routing processor connected to said network interface via an intra-device communication line and performing routing of a packet received from said network interface;

wherein said routing processor includes:

path information holding means; and

next-path searching means for calculating, based on path information held in said path information holding means, a path to which the received packet is to be forwarded next

wherein said next-path searching means searches for a next path using a 2-branch tree search in which a destination address is checked, one bit at a

time, beginning with a highest-order bit, stores 0- to k-bit mask 2-branch tree nodes in internal storing means of said next-path searching means, a number of bits of each of said nodes being equal to or smaller than a predetermined number, stores (k+1)bit or longer mask 2-branch tree nodes in external storing means of said search means, and performs pipeline processing for 0- to k-bit mask node search processing and (k+1)bit or longer mask node search processing.

10. A network next-hop search method for use in a network forwarding device connected to a plurality of network and transmitting a packet received from one of said networks to a next hop based on path information

wherein, in a 2-branch tree search in which a destination address of the received packet is checked, one bit at a time beginning at a highest-order bit, p (p is an integer equal to or larger than 2) levels are combined into one 2^p -branch tree to perform a search of the p levels of the 2-branch tree as a one-level search.

ABSTRACT

The present invention improves a network forwarding device, such as a router, in a computer network system. When an address to which a received packet is to be forwarded is searched for based on path information, a path search data structure is changed to the one in which, in a 2-branch tree search in which a destination address of the received packet is checked, one bit at a time beginning at a highest-order bit, p (p is an integer equal to or larger than 2) levels are combined into one 2^p -branch tree to perform a search of the p levels of the 2-branch tree as a one-level search. This data structure speeds up path search processing.

FIG. 1

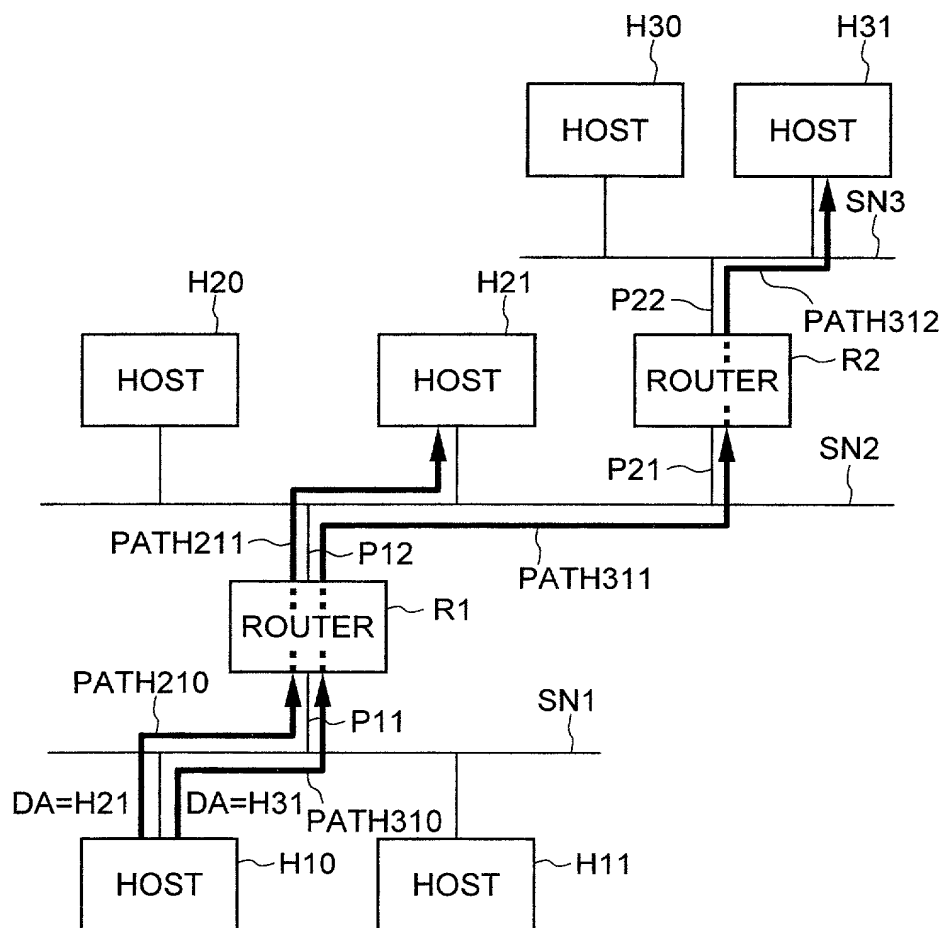


FIG. 2

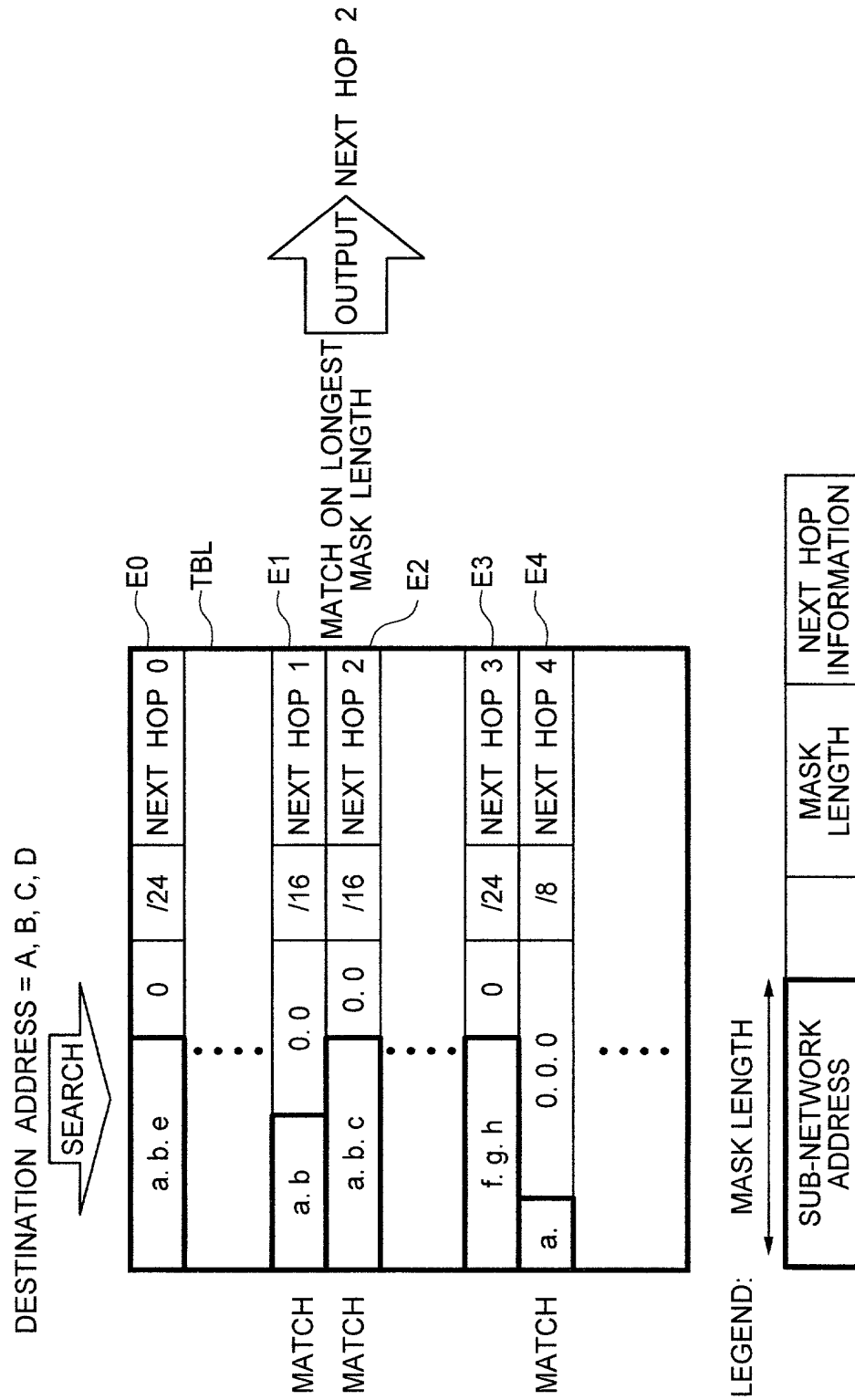
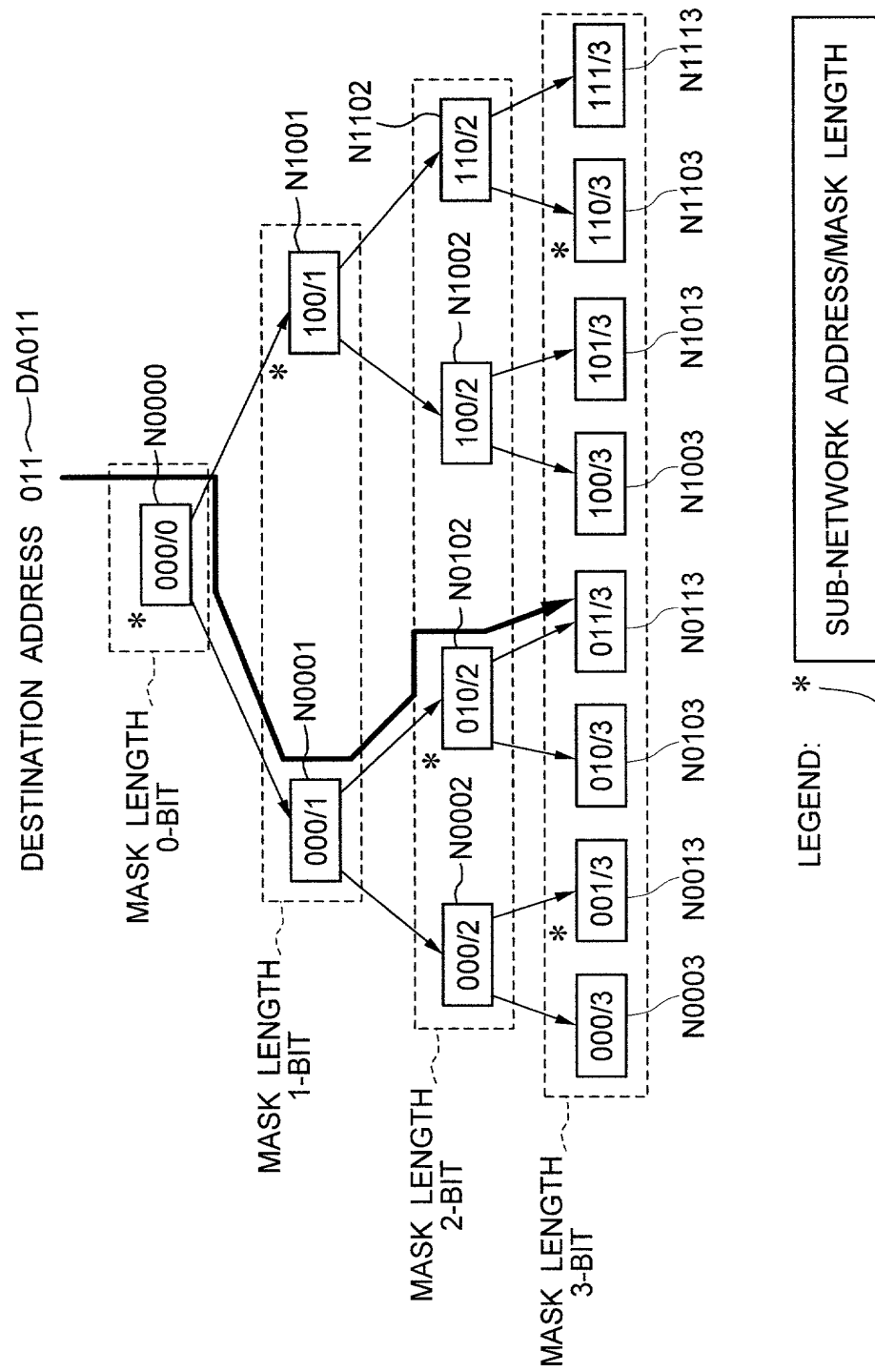


FIG. 3



LEGEND:

SUB-NETWORK ADDRESS/MASK LENGTH

INDICATES THAT A PATH IS ALLOCATED TO THIS NODE.

FIG. 4

PATH TABLE

SUB-NETWORK ADDRESS	/	MASK LENGTH
000	/	0
001	/	3
010	/	2
100	/	1
110	/	3

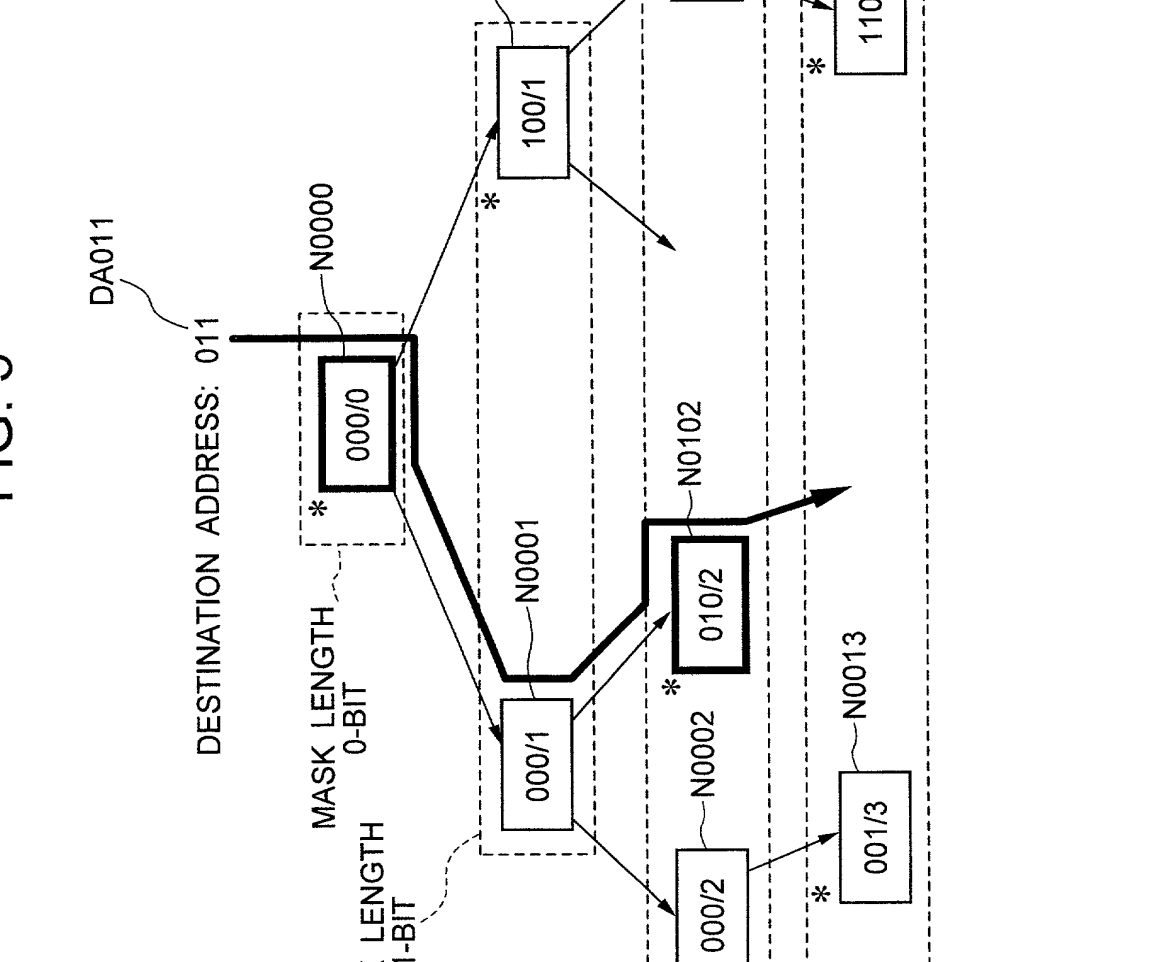


FIG. 6

PATH TABLE

SUB-NETWORK ADDRESS	/	MASK LENGTH
0x00. 00. 00. 00	/	0
0x85. 04. 00. 00	/	16
0x85. 05. 00. 00	/	16

7/30

FIG. 7

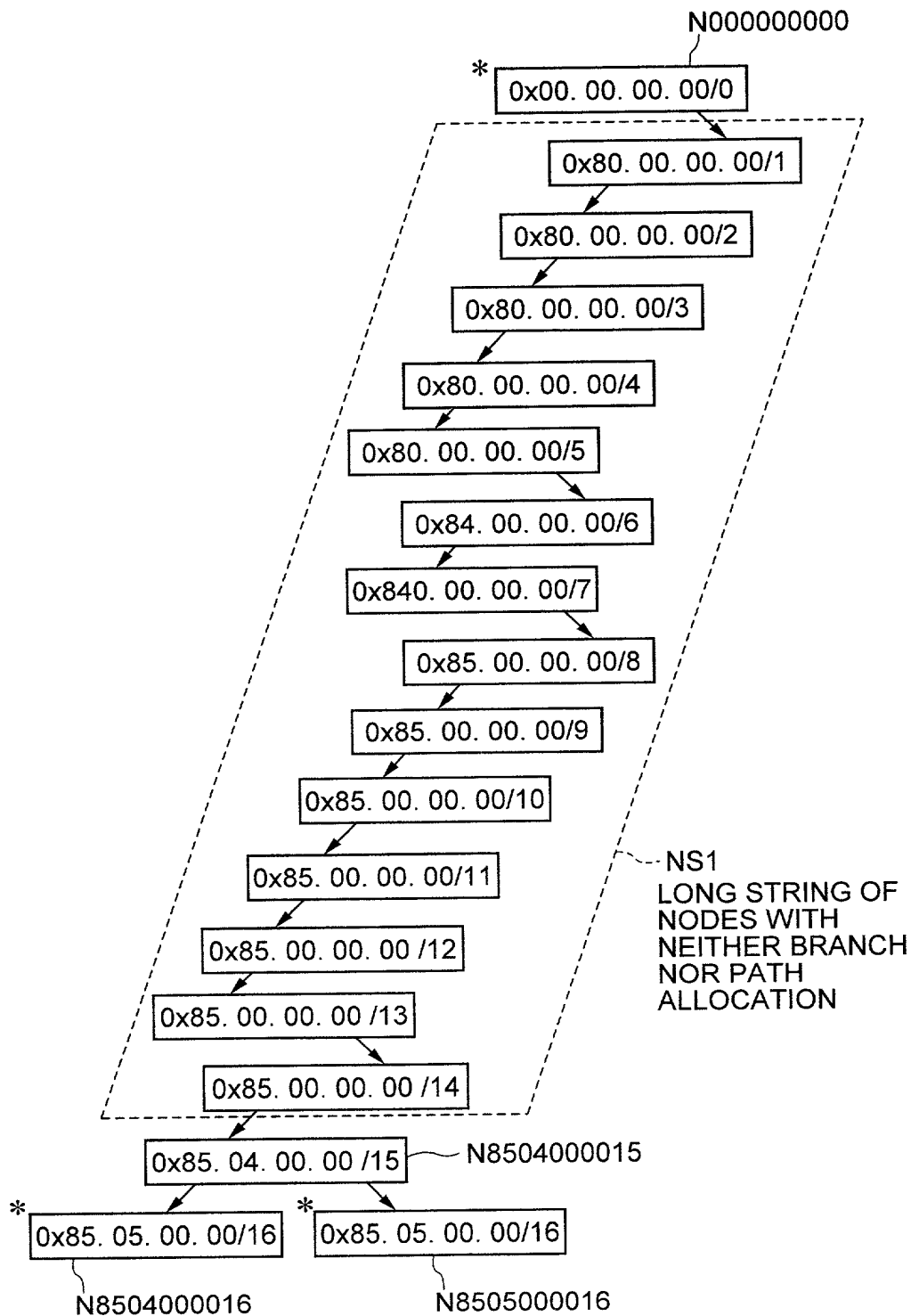
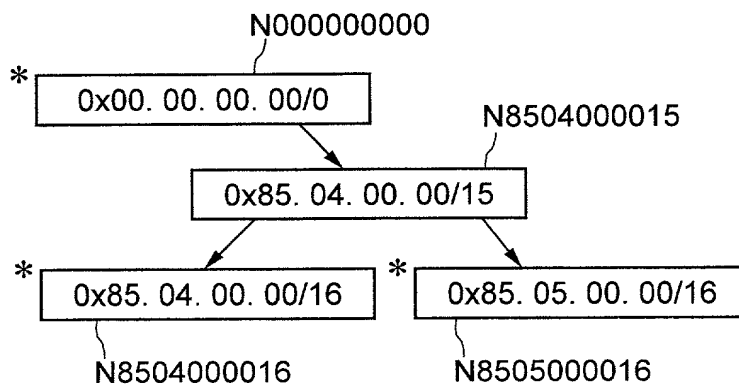


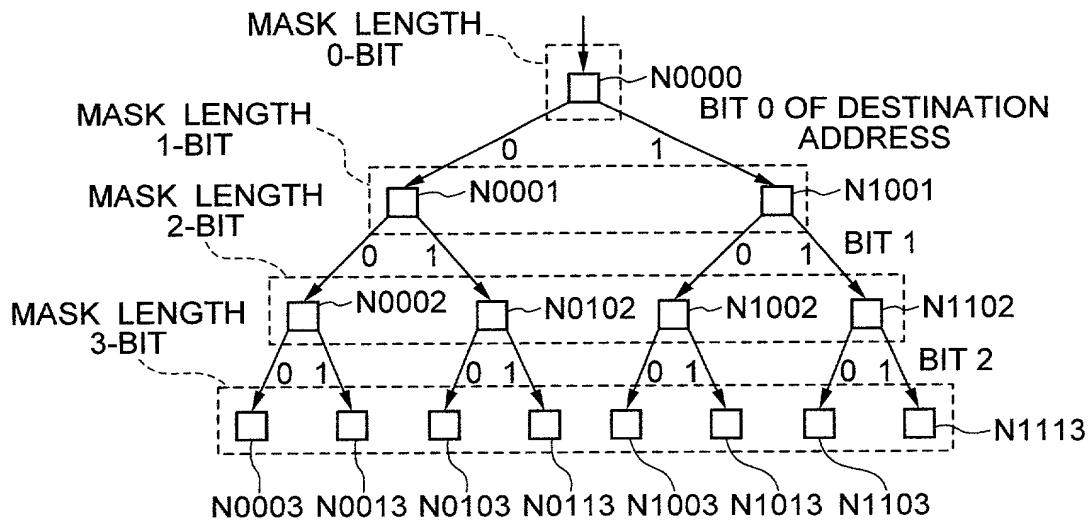
FIG. 8



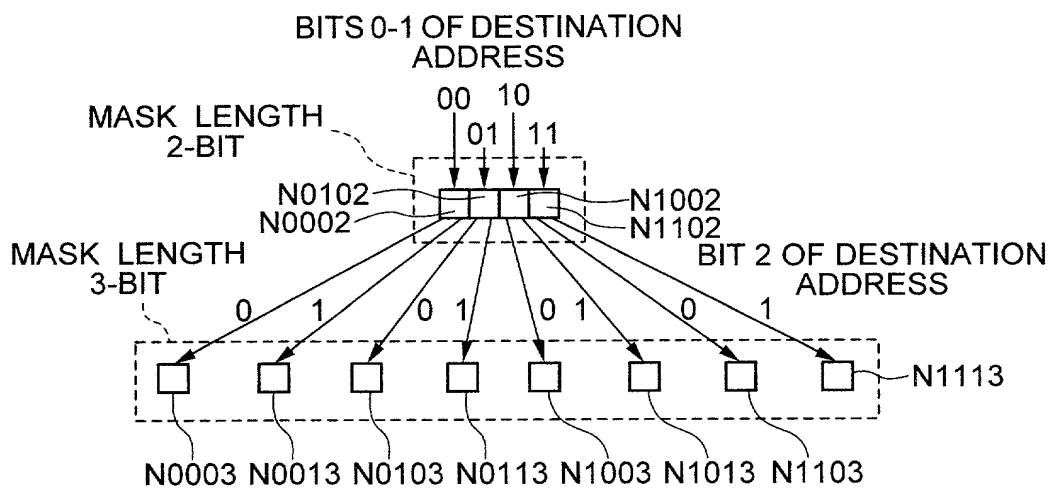
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FIG. 9

(a)



(b)



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FIG. 10

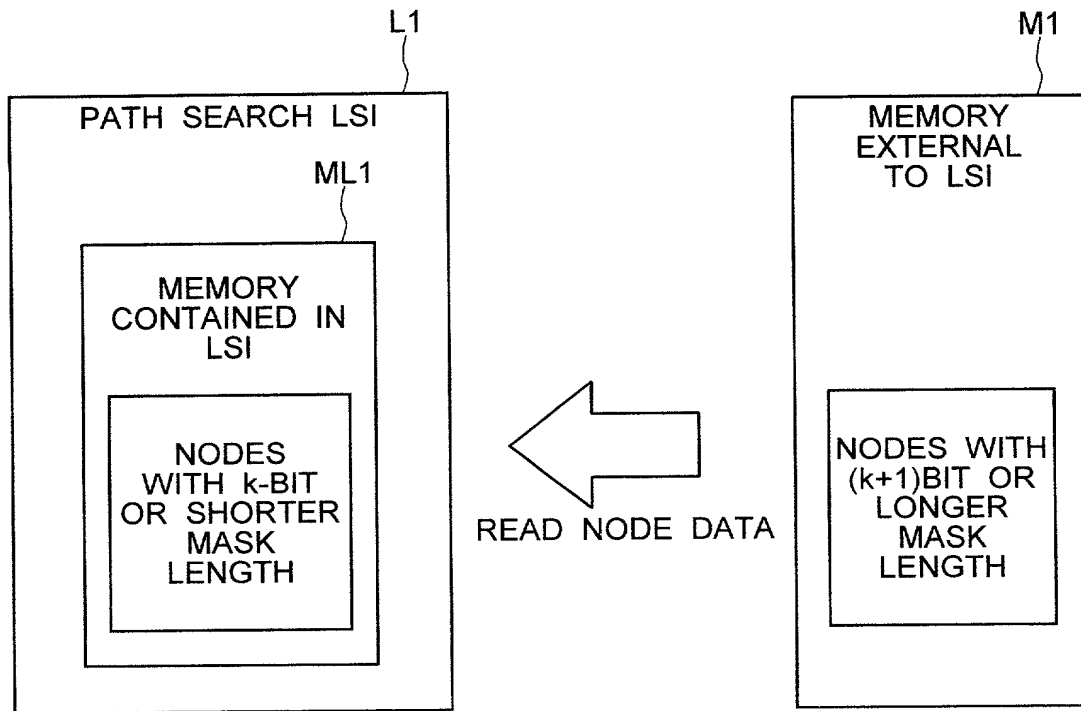


FIG. 11

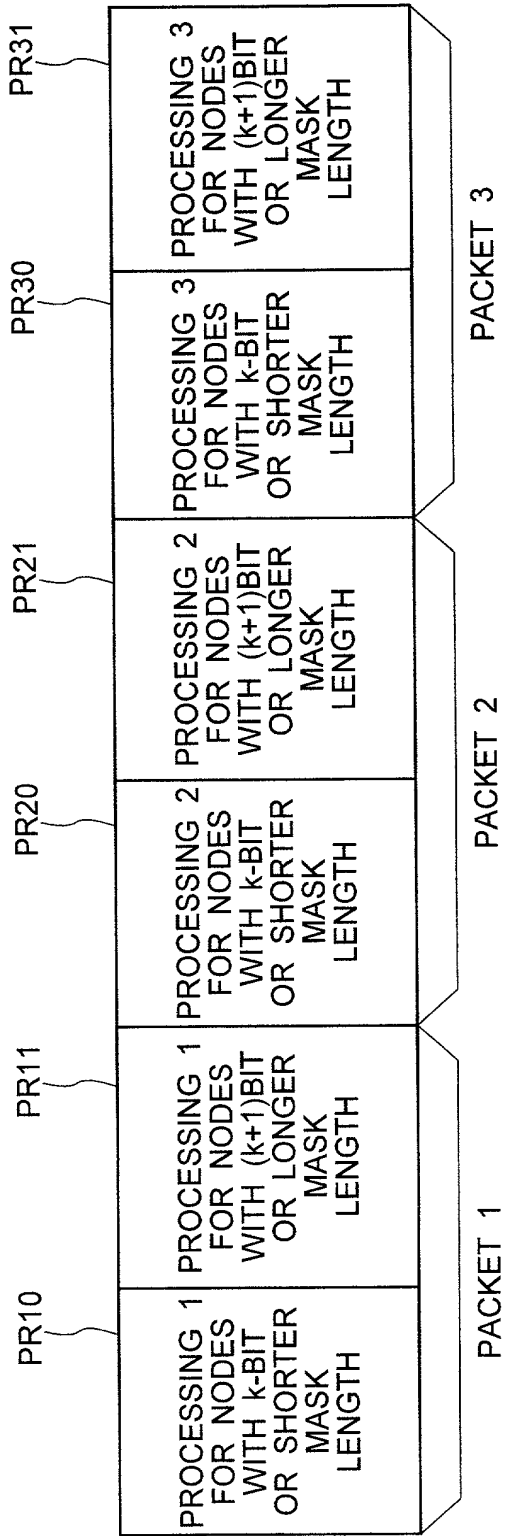


FIG. 12

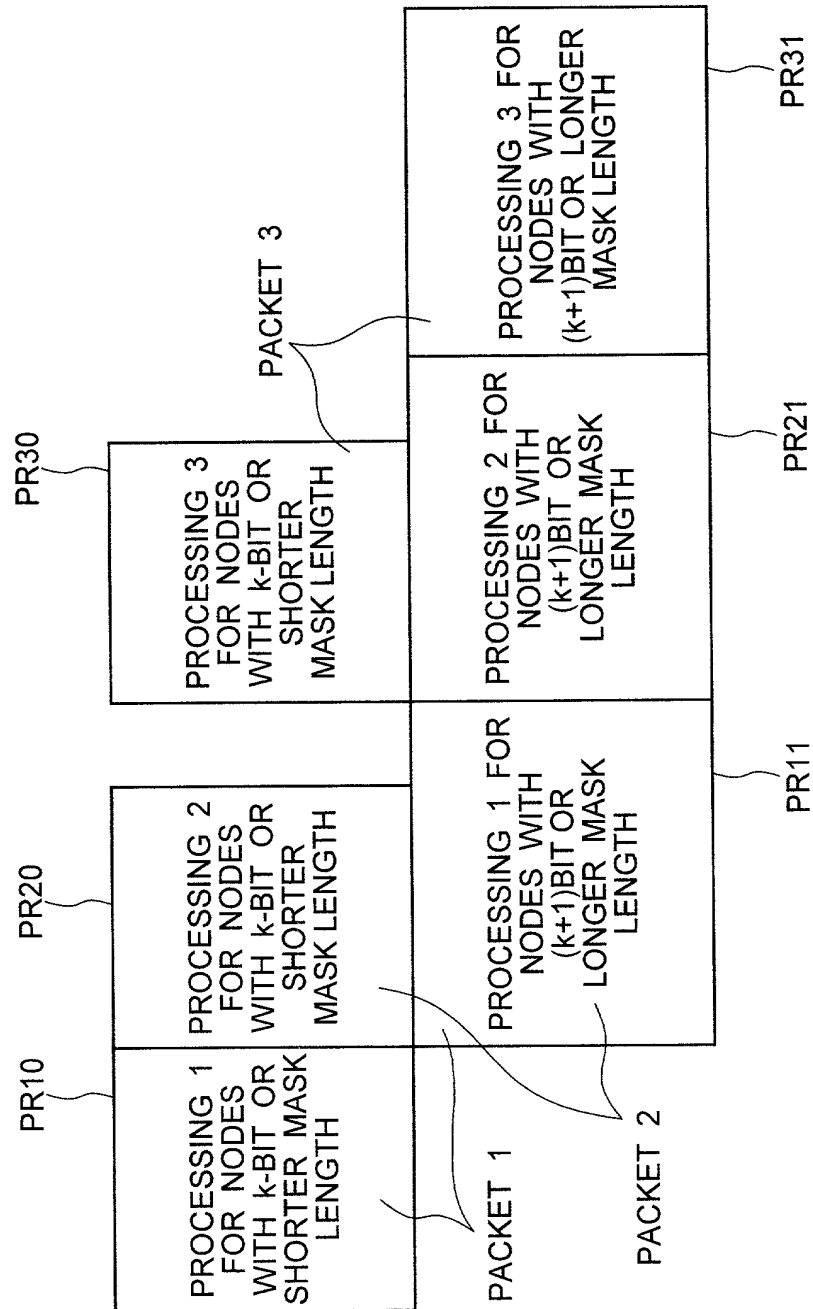


FIG. 13

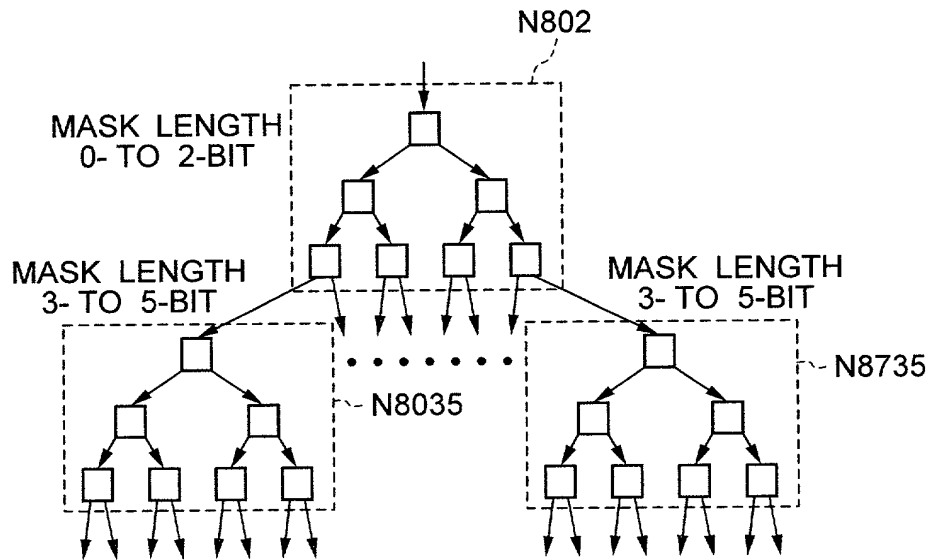


FIG. 14

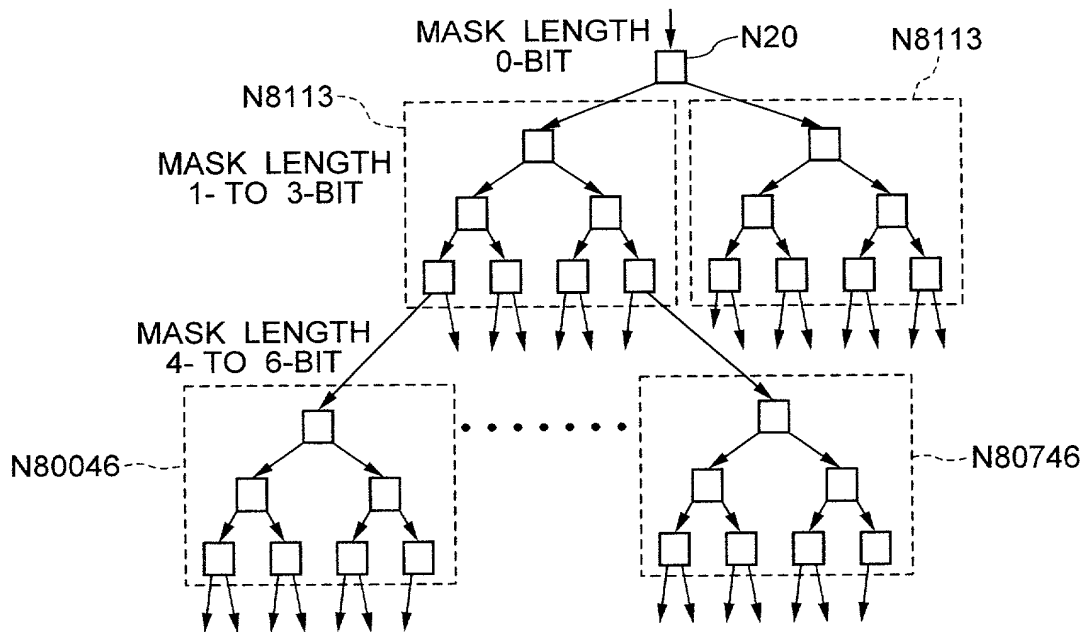
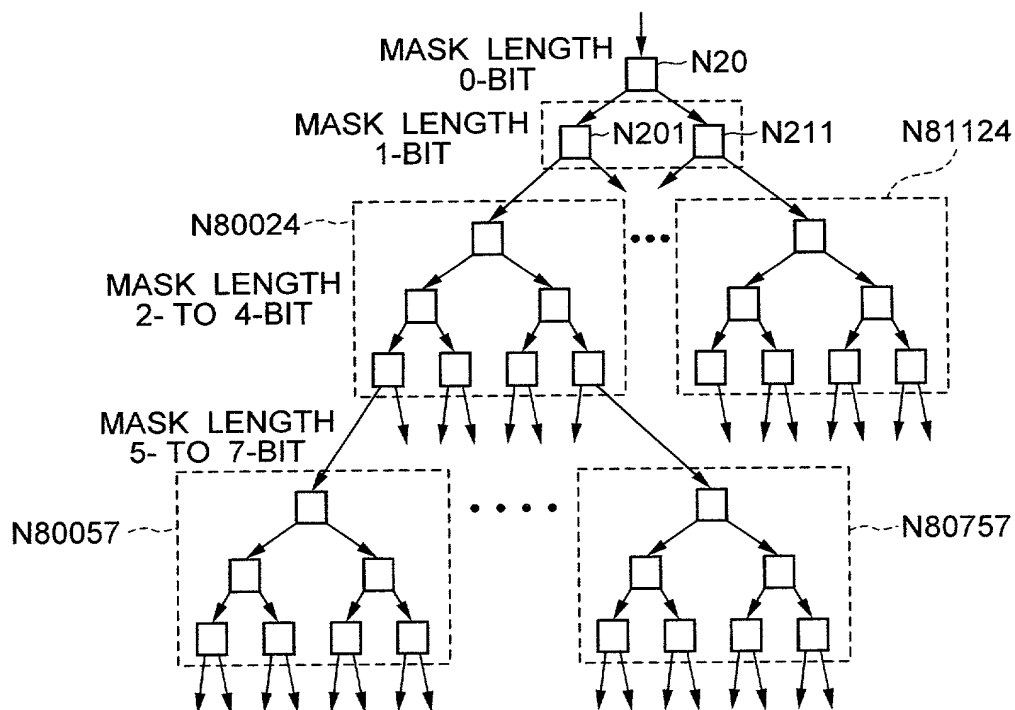


FIG. 15



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FIG. 16

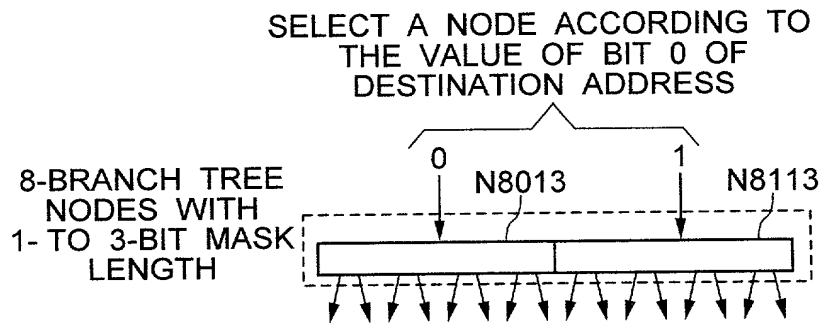


FIG. 17

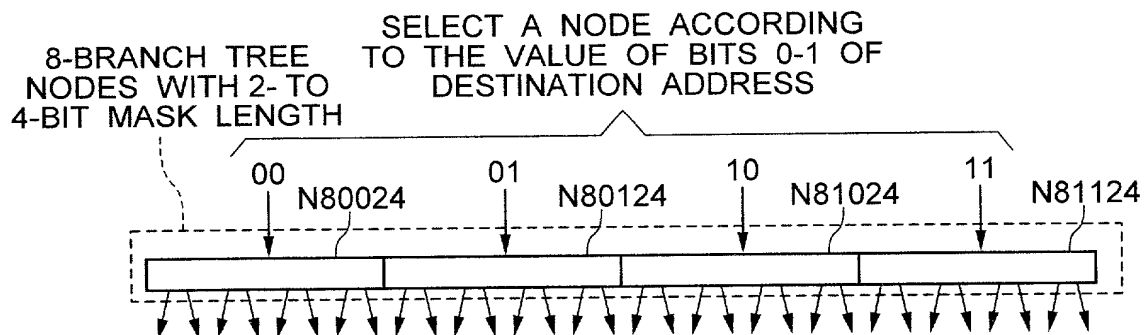


FIG. 18

SELECT A NODE ACCORDING
TO THE VALUE OF BITS 0-2 OF
DESTINATION ADDRESS

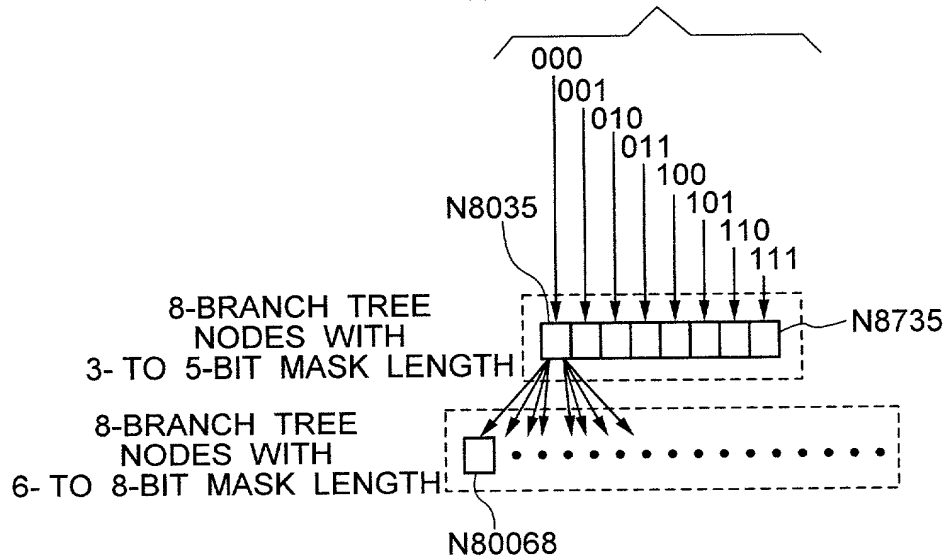


FIG. 19

SELECT A NODE ACCORDING
TO THE VALUE OF BITS 0-3 OF
DESTINATION ADDRESS

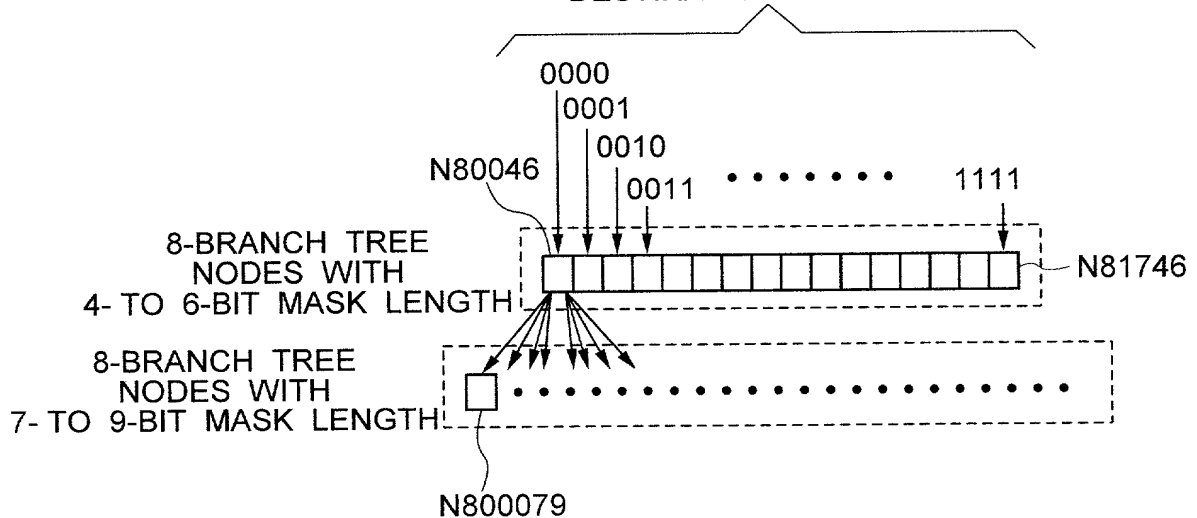
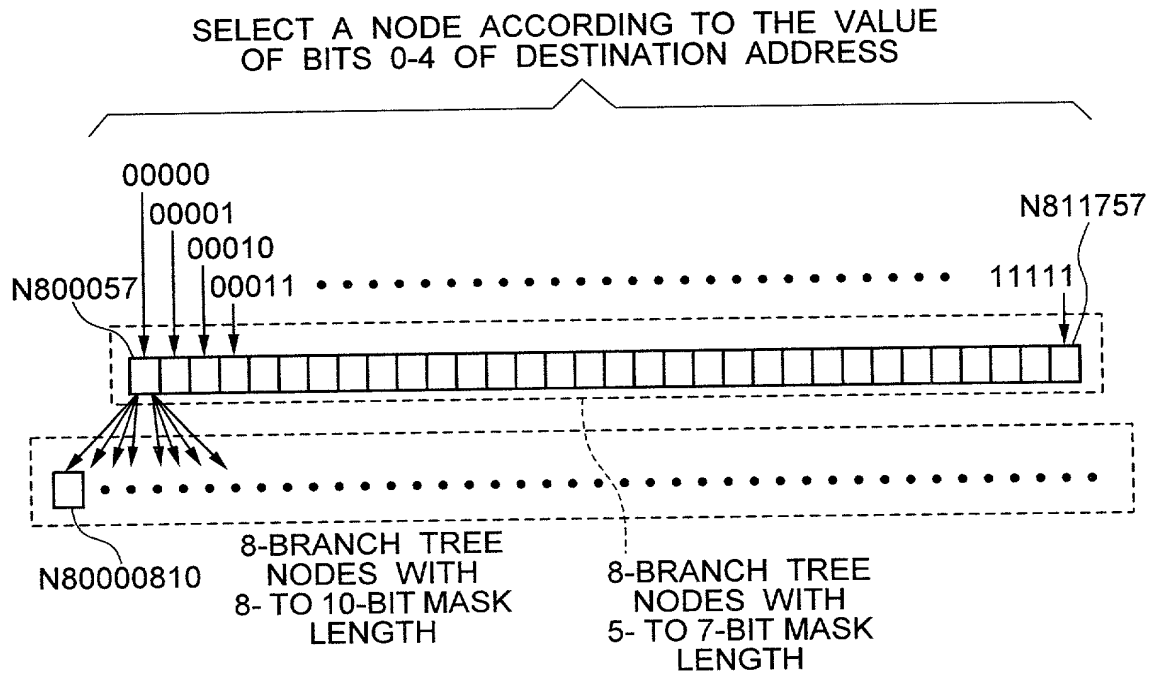
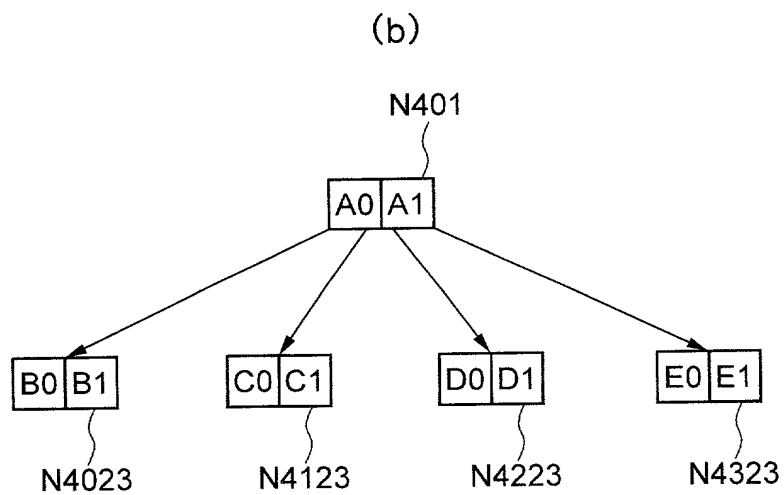
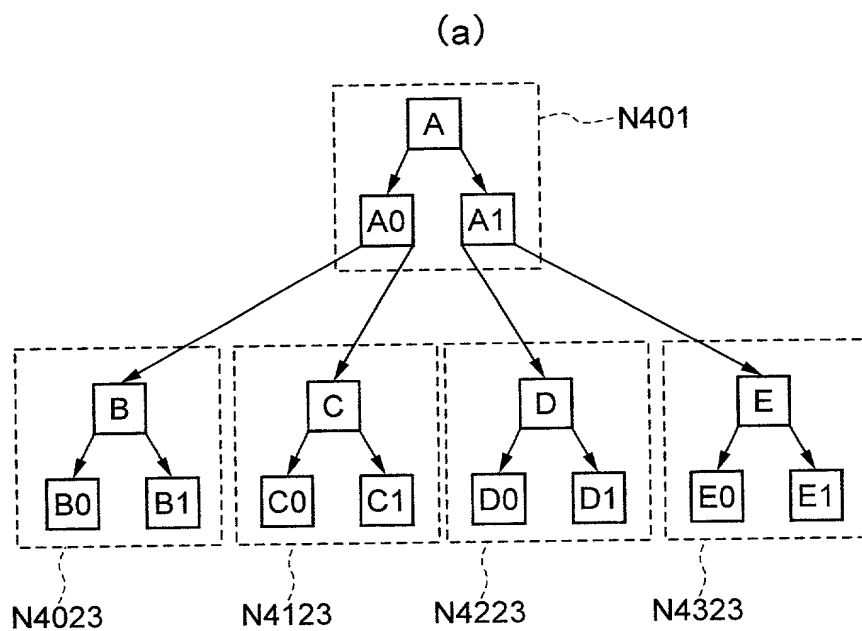


FIG. 20



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FIG. 21



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FIG. 22

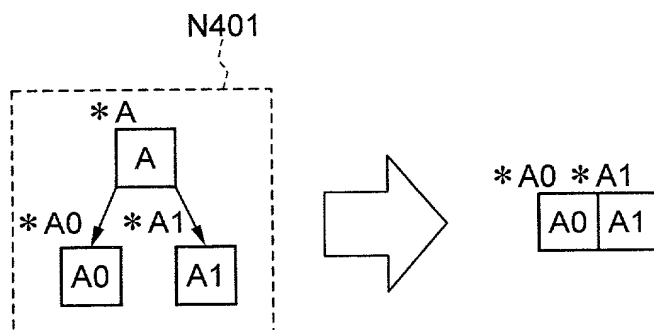


FIG. 23

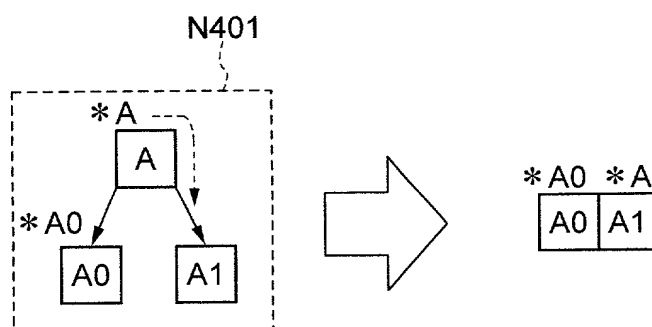
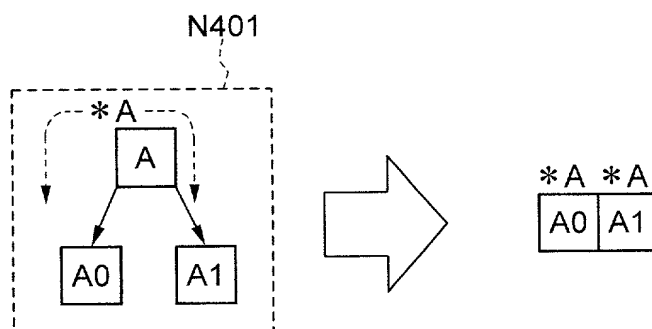


FIG. 24



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FIG. 25

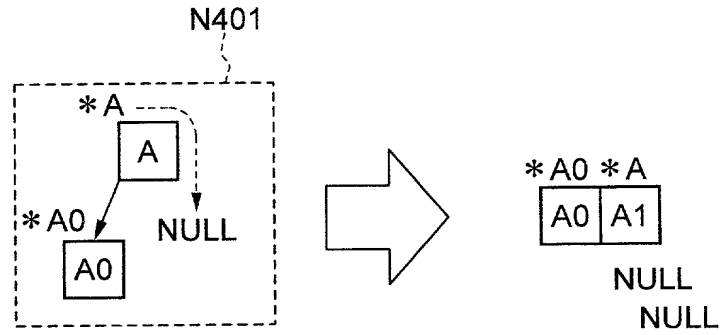


FIG. 26

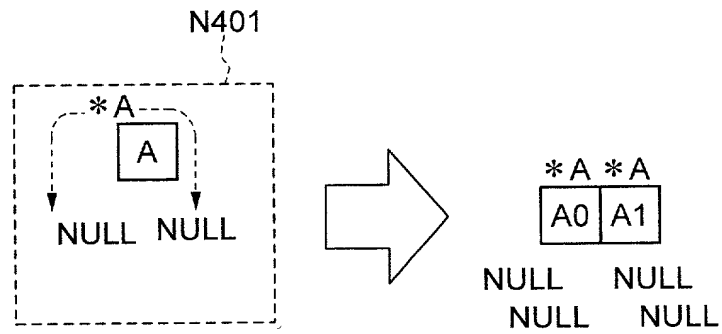


FIG. 27

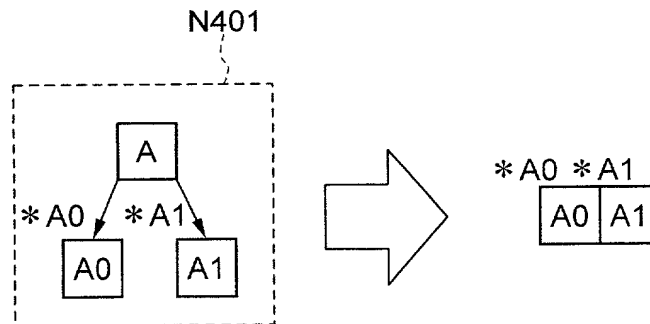


FIG. 28

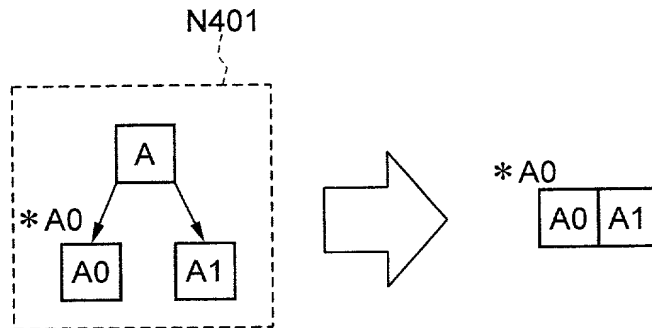


FIG. 29

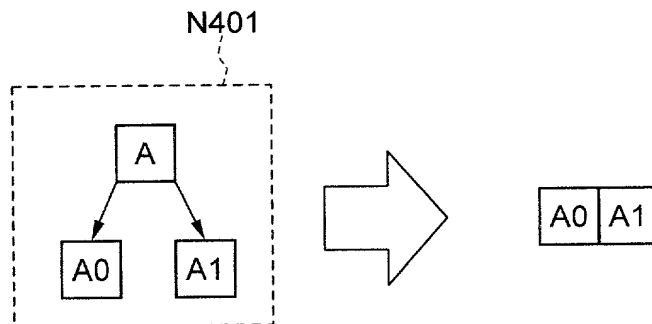


FIG. 30

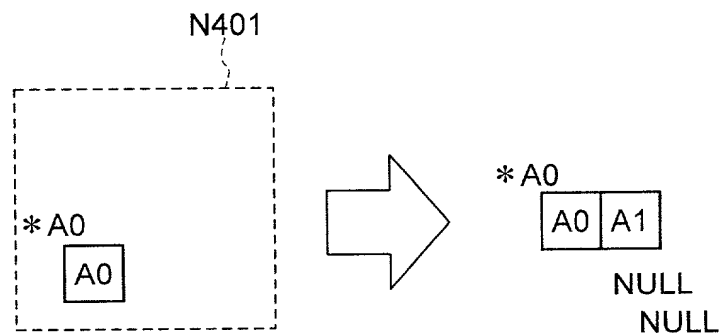
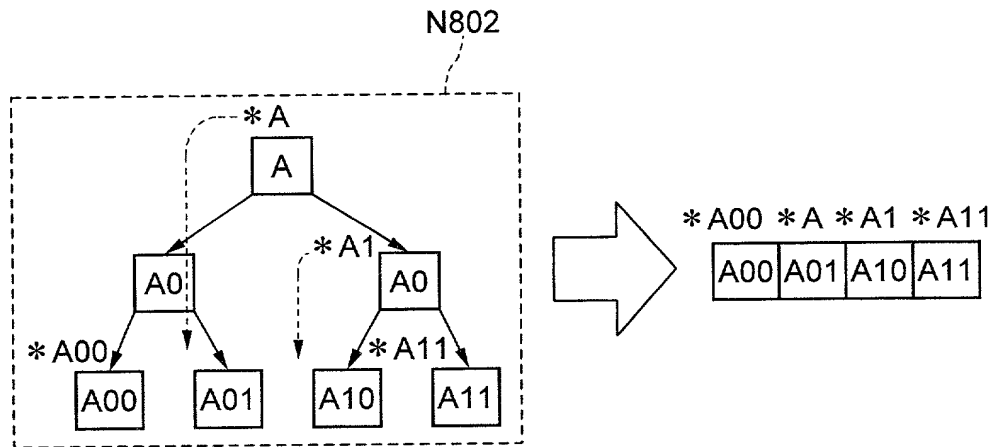


FIG. 31

(a)



(b)

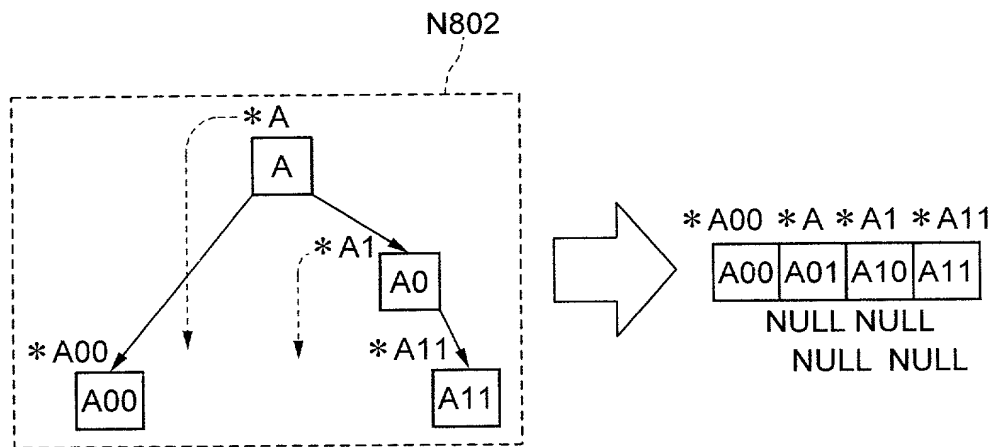


FIG. 32

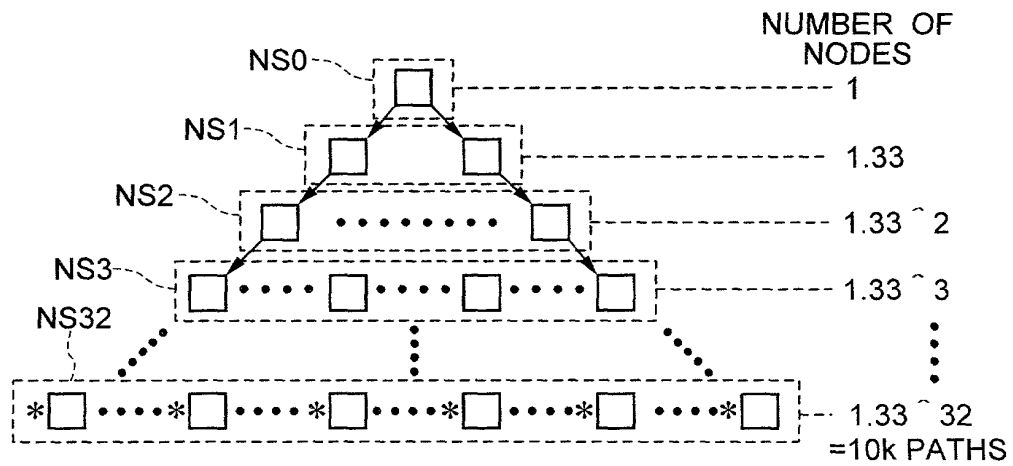
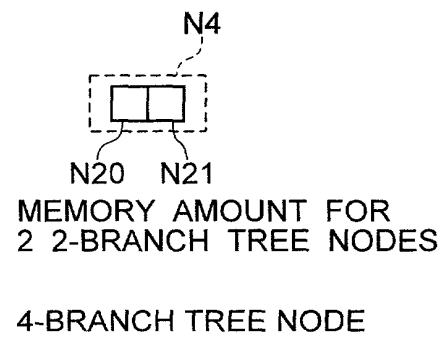
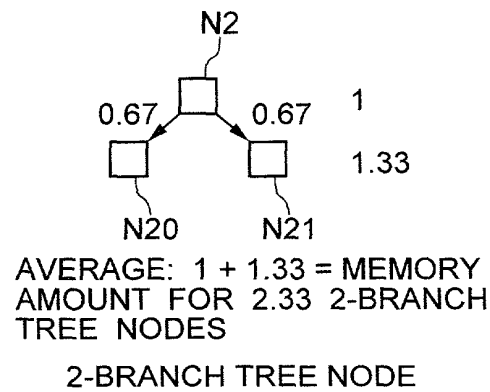


FIG. 33

(a)



(b)

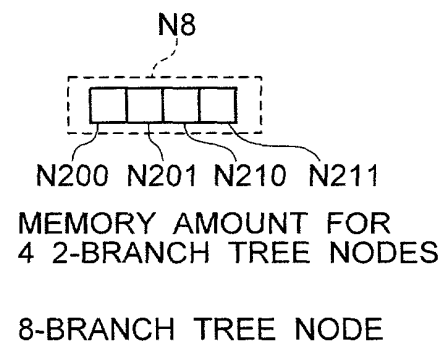
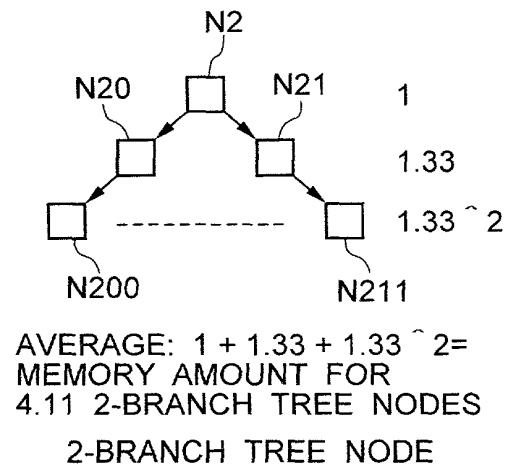
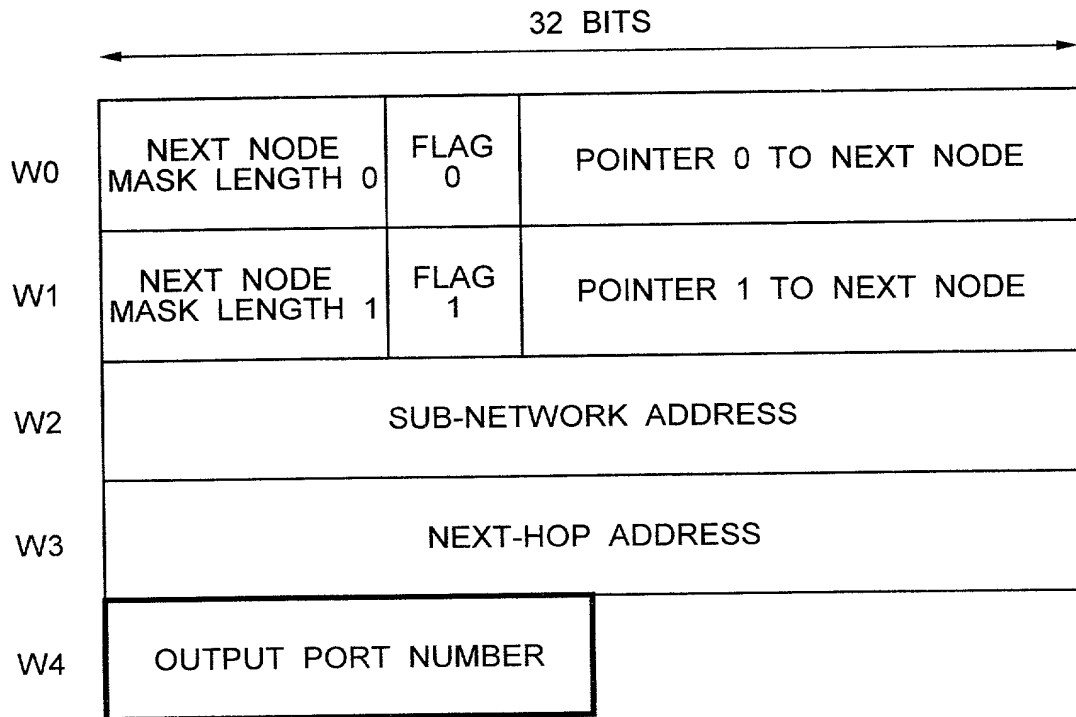
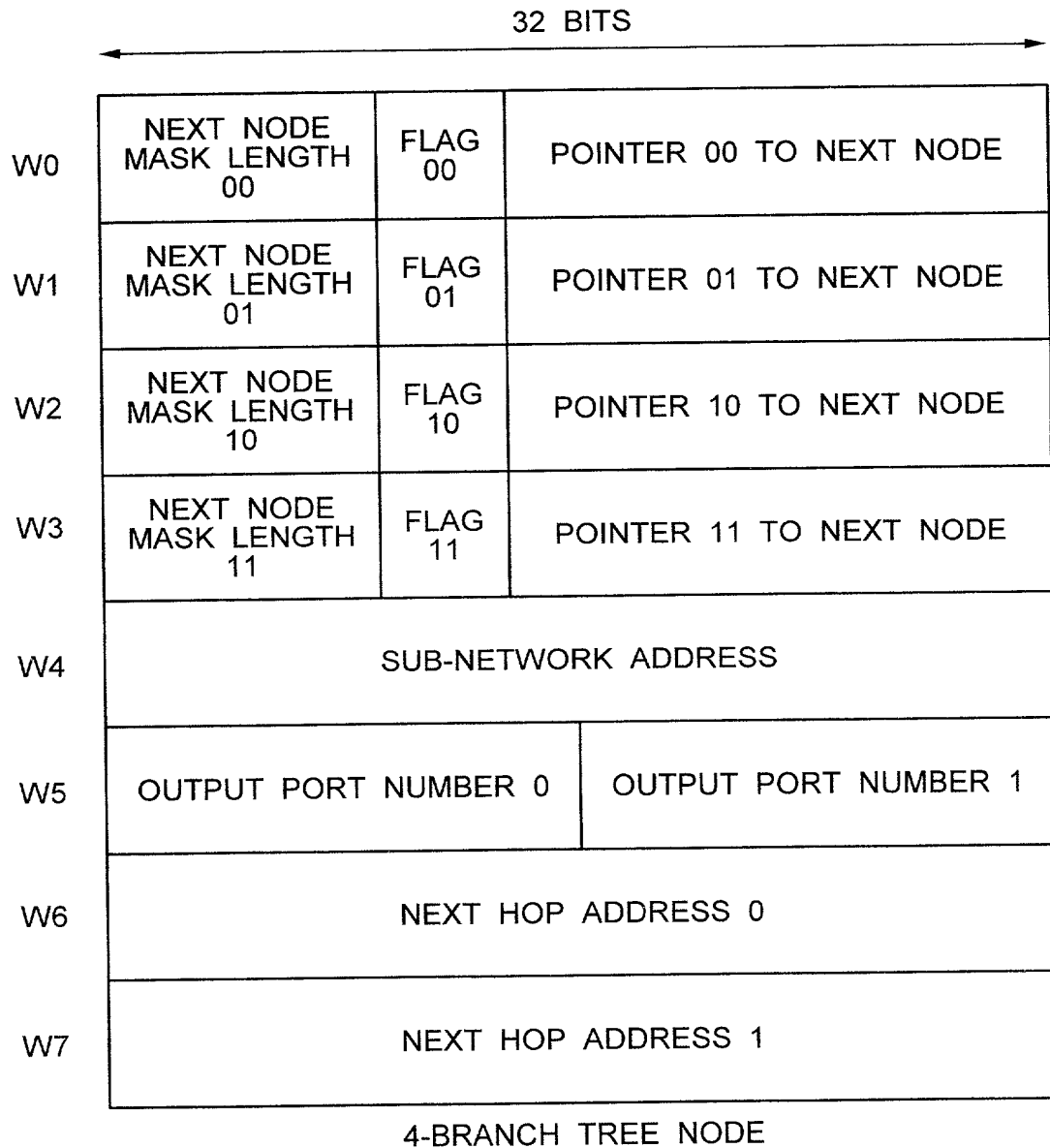


FIG. 34



2-BRANCH TREE NODE

FIG. 35



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FIG. 36

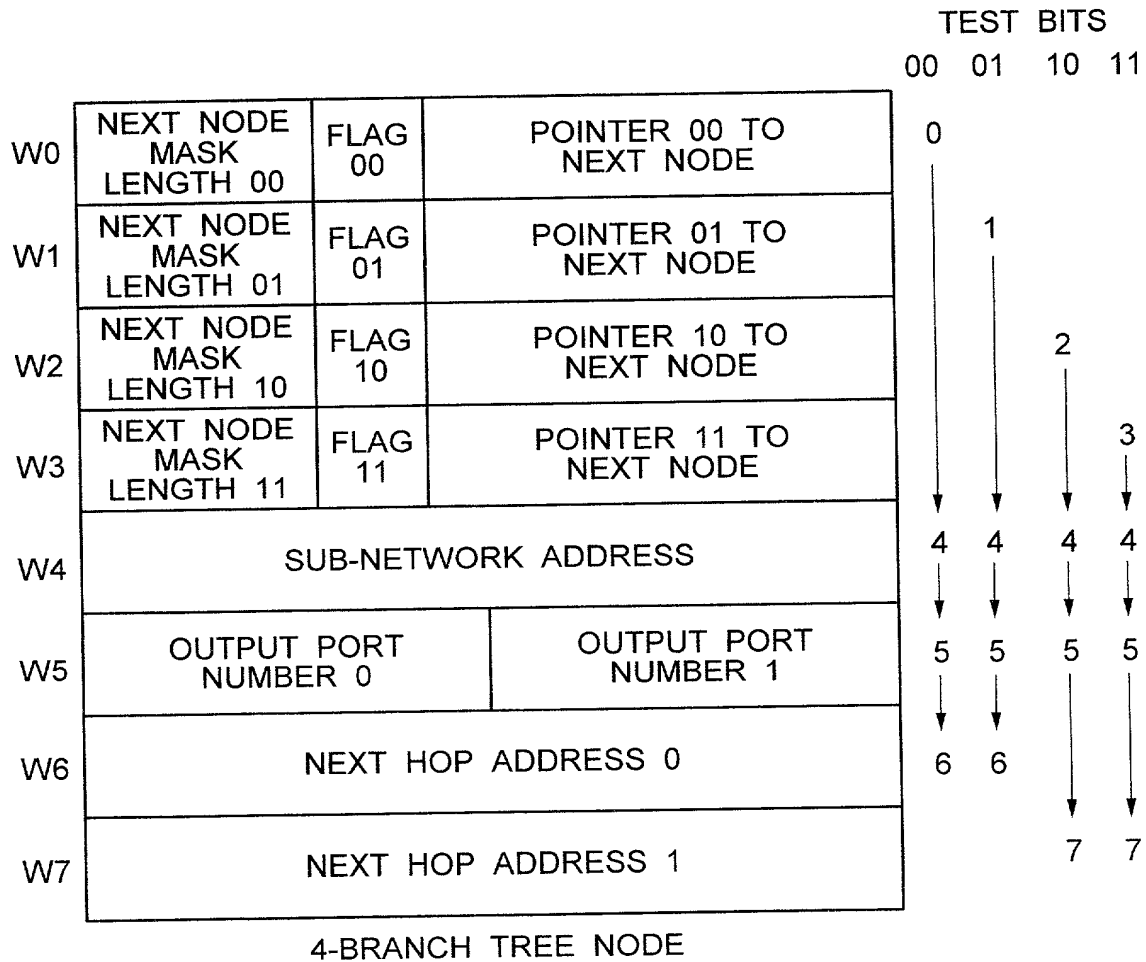


FIG. 37

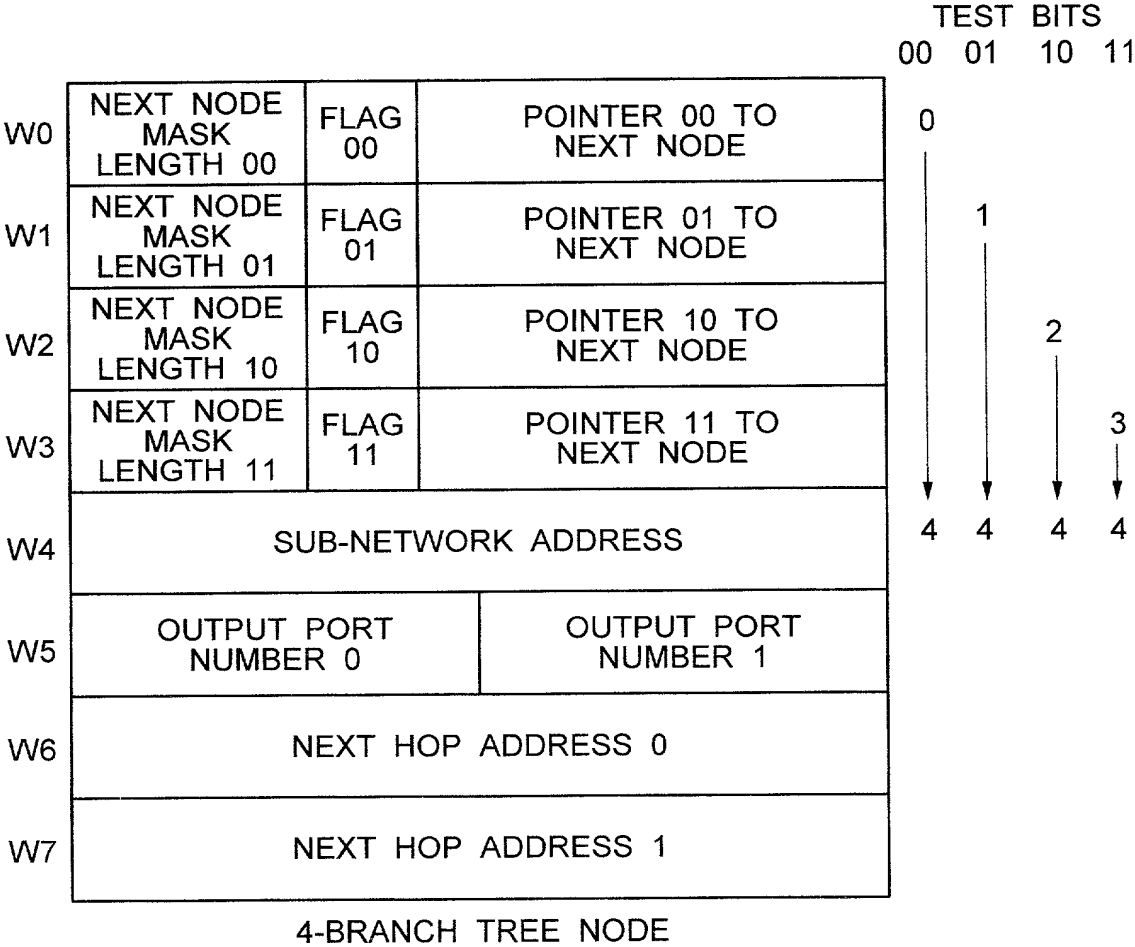


FIG. 38

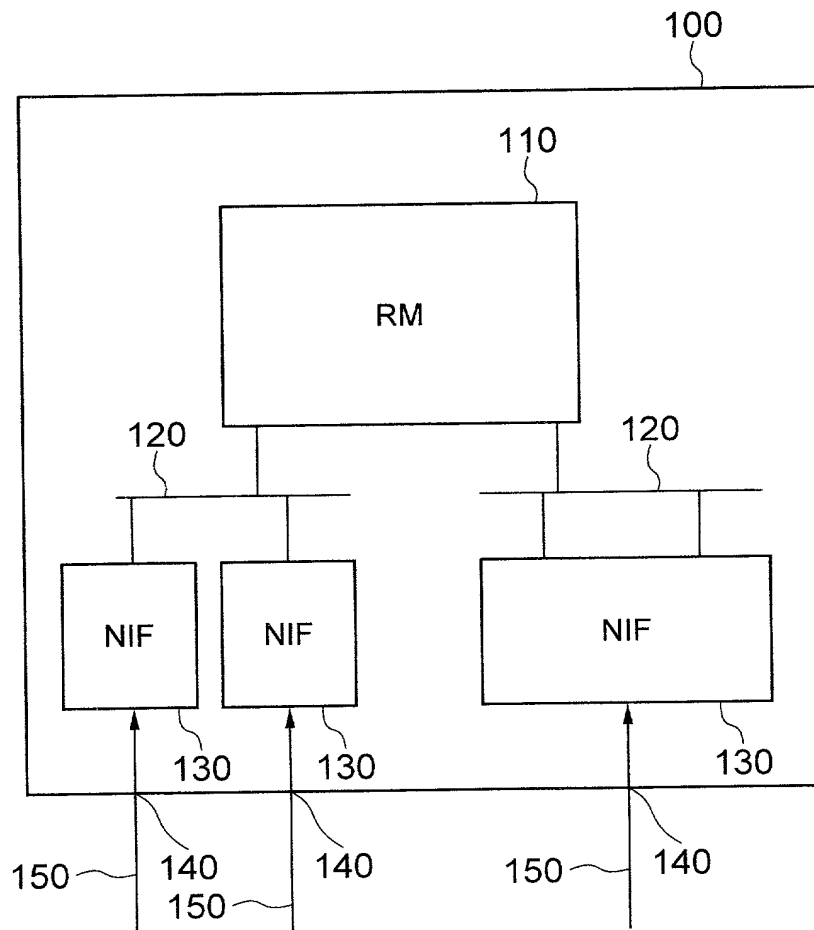
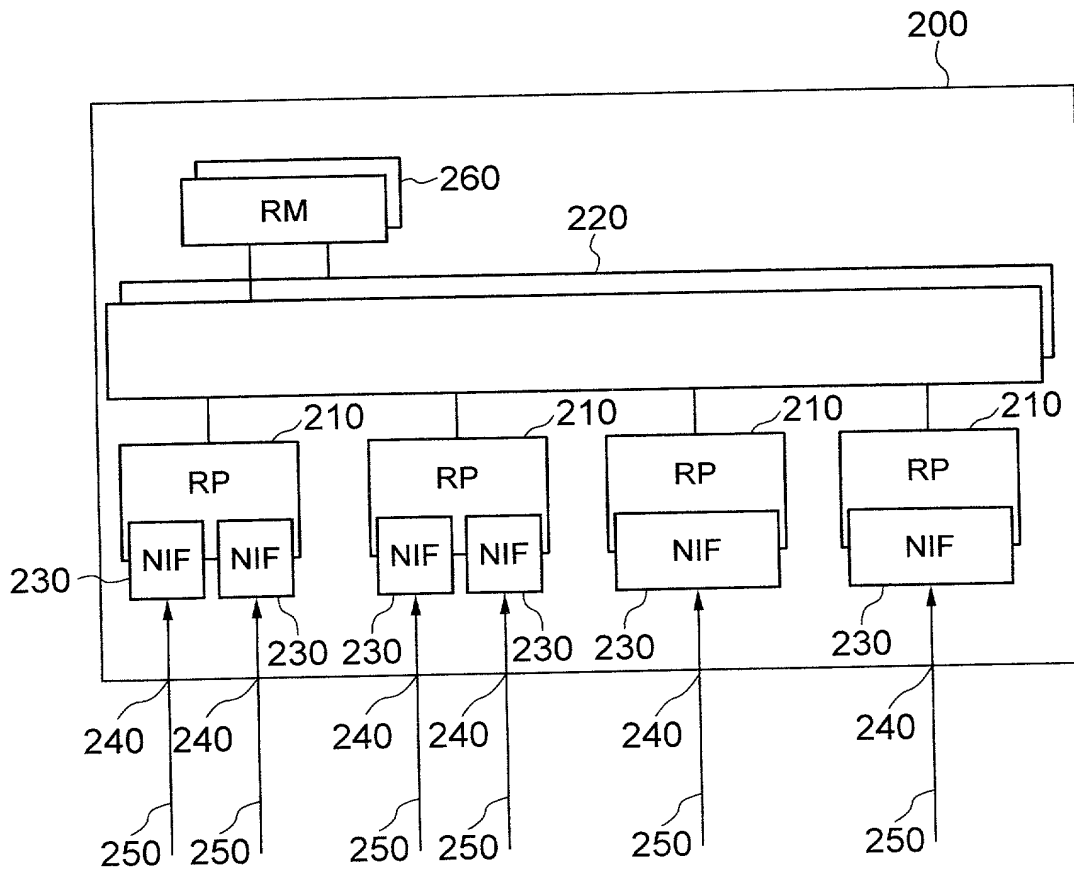


FIG. 39



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(*)

PTO/SB/106(8-96)

Approved for use through 9/30/98. OMB 0651-0032

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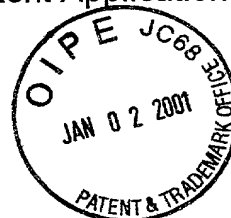
Under the Paperwork Reduction Act of 1995, no persons are required to respond to a collection of information unless it displays a valid OMB control number.

Declaration and Power of Attorney For Patent Application

特許出願宣言書及び委任状

Japanese Language Declaration

日本語宣言書



下記の氏名の発明者として、私は以下の通り宣言します。

As a below named inventor, I hereby declare that:

私の住所、私書箱、国籍は下記の私の氏名の後に記載された通りです。

My residence, post office address and citizenship are as stated next to my name.

下記の名称の発明に関して請求範囲に記載され、特許出願している発明内容について、私が最初かつ唯一の発明者（下記の氏名が一つの場合）もしくは最初かつ共同発明者であると（下記の名称が複数の場合）信じています。

I believe I am the original, first and sole inventor (if only one name is listed below) or an original, first and joint inventor (if plural names are listed below) of the subject matter which is claimed and for which a patent is sought on the invention entitled

NETWORK FORWARDING DEVICE AND NETWORK

NEXT-HOP SEARCH METHOD

上記発明の明細書（下記の欄で×印がついていない場合は、本書に添付）は、

The specification of which is attached hereto unless the following box is checked:

☐ __月__日に提出され、米国出願番号または特許協定条約国際出願番号を____とし、
(該当する場合) _____に訂正されました。

☒ was filed on August 17, 2000
as United States Application Number or
PCT International Application Number
09/622,484 and was amended on
_____(if applicable).

私は、特許請求範囲を含む上記訂正後の明細書を検討し、内容を理解していることをここに表明します。

I hereby state that I have reviewed and understand the contents of the above identified specification, including the claims, as amended by any amendment referred to above.

私は、連邦規則法典第37編第1条56項に定義されるとおり、特許資格の有無について重要な情報を開示する義務があることを認めます。

I acknowledge the duty to disclose information which is material to patentability as defined in Title 37, Code of Federal Regulations, Section 1.56.

Page 1 of 5

Burden Hour Statement This form is estimated to take 0.4 hours to complete. Time will vary depending upon the needs of the individual case. Any comments on the amount of time you are required to complete this form should be sent to the Chief Information Officer, Patent and Trademark Office, Washington, DC 20231. DO NOT SEND FEES OR COMPLETED FORMS TO THIS ADDRESS. SEND TO Commissioner of Patents and Trademarks, Washington, DC 20231.

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Japanese Language Declaration
(日本語宣言書)

私は、米国法典第35編119条(a)-(d)項又は365条(b)項に基づき下記の、米国以外の国の少なくとも一カ国を指定している特許協力条約365(a)項に基づく国際出願、又は外国での特許出願もしくは発明者証の出願についての外国優先権をここに主張するとともに、優先権を主張している、本出願の前に出願された特許または発明者証の外国出願を以下に、枠内をマークすることで、示している。

Prior Foreign Application(s): NONE

外国での先行出願

Priority Not Claimed

優先権主張なし

(Number) (番号)	(Country) (国名)	(Day/Month/Year Filed) (出願年月日)
(Number) (番号)	(Country) (国名)	(Day/Month/Year Filed) (出願年月日)

私は、第35編米国法典119条(e)項に基づいて下記の米国特許出願規定に記載された権利をここに主張いたします。

I hereby claim the benefit under Title 35, United States Code, Section 119(e) of any United States provisional application(s) listed below.

(Application No.) (出願番号)	(Filing Date) (出願日)	(Application No.) (出願番号)	(Filing Date) (出願日)
-----------------------------	------------------------	-----------------------------	------------------------

私は、下記の米国法典第35編120条に基づいて下記の米国特許出願に記載された権利、又は米国を指定している特許協力条約365条(c)に基づく権利をここに主張します。また、本出願の各請求範囲の内容が米国法典第35編112条第1項又は特許協力条約で規定された方法で先行する米国特許出願に開示されていない限り、その先行米国出願書提出日以降で本出願書の日本国内または特許協力条約国際提出日までの期間中に入手された、連邦規則法典第37編1条56項で定義された特許資格の有無に関する重要な情報について開示義務があることを認識しています。

I hereby claim the benefit under Title 35, United States Code, Section 120 of any United States application(s), or 365(c) of any PCT international application designating the United States, listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in the prior United States or PCT International application in the manner provided by the first paragraph of Title 35, United States Code Section 112, I acknowledge the duty to disclose information which is material to patentability as defined in Title 37, Code of Federal Regulations, Section 1.56 which became available between the filing date of the prior application and the national or PCT international filing date of application.

PCT/JP98/01232	23/March/1998	Pending
(Application No.) (出願番号)	(Filing Date) (出願日)	(Status: Patented, Pending, Abandoned) (現況: 特許許可済、係属中、放棄済)
(Application No.) (出願番号)	(Filing Date) (出願日)	(Status: Patented, Pending, Abandoned) (現況: 特許許可済、係属中、放棄済)

私は、私自身の知識に基づいて本宣言書で私が行なう表明が真実であり、かつ私の入手した情報と私の信じることに基き、かつ表明が全て真実であると信じていること、さらに故意になされた虚偽の表明及びそれと同等の行為は米国法典第18編第1001条に基づき、罰金または拘禁、もしくはその両方により処罰されること、そしてそのような故意による虚偽の表明を行なえば、出願した、又は既に許可された特許の有効性が失われることを認識し、よってここに上記のごとく宣誓を致します。

I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code and that such willful false statements may jeopardize the validity of the application or any patent issued thereon.

Japanese Language Declaration (日本語宣言書)

委任状： 私は下記の発明者として、本出願に関する一切の手続きを米特許商標局に対して遂行する弁理士または代理人として、下記の者を指名いたします。(弁理士、または代理人の氏名及び登録番号を明記のこと)

POWER OF ATTORNEY: As a named inventor, I hereby appoint the following attorney(s) and/or agent(s) to prosecute this application and transact all business in the Patent and Trademark Office connected therewith (list name and registration number)

Donald R. Antonelli, Reg. No. 20,296; David T. Terry, Reg. No. 20,178; Melvin Kraus, Reg. No. 22,466; William I. Solomon, Reg. No. 28,565; Gregory E. Montone, Reg. No. 28,141; Ronald J. Shore, Reg. No. 28,577; Donald E. Stout, Reg. No. 26,422; Alan E. Schiavelli, Reg. No. 32,087; James N. Dresser, Reg. No. 22,973 and Carl I. Brundidge, Reg. No. 29,621

書類送付先

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Suite 1800
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Arlington, Virginia 22209

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Direct Telephone Calls to: (name and telephone number)

Telephone: (703) 312-6600
Fax: (703) 312-6666

唯一または第一発明者	Full name of sole or first inventor	
	Kazuo SUGAI	
発明者の署名	日付	Inventor's signature Date
		Kazuo Sugai 10/19/00
住所	Residence	
	Zama, Japan JAX	
国籍	Citizenship	
	Japan	
私書箱	Post Office Address	
	c/o Hitachi, Ltd., Intellectual Property Group New Marunouchi Bldg. 5-1, Marunouchi 1-chome, Chiyoda-ku, Tokyo 100-8220, Japan	

(第二以降の共同発明者についても同様に記載し、署名をすること)

(Supply similar information and signature for second and subsequent joint inventors.)

第二共同発明者		2nd		Full name of second joint inventor, if any Takeshi AIMOTO	
第二共同発明者の署名	日付	Takeshi Aimoto		Second inventor's signature	Date 12/06/00
住所	Residence Sagamihara, Japan JPX				
国籍	Citizenship Japan				
私書箱	Post Office Address c/o Hitachi, Ltd., Intellectual Property Group New Marunouchi Bldg. 5-1, Marunouchi 1-chome, Chiyoda-ku, Tokyo 100-8220, Japan				
第三共同発明者		3rd		Full name of third joint inventor, if any Nobuhito MATSUYAMA	
第三共同発明者の署名	日付	Nobuhito Matsuyama		Third inventor's signature	Date 10/30/00
住所	Residence Hadano, Japan JPX				
国籍	Citizenship Japan				
私書箱	Post Office Address c/o Hitachi Information Technology Co., Ltd., 456, Sakai, Nakai-machi, Ashigarakami-gun, Kanagawa-ken, 259-0157, Japan				
第四共同発明者		4th		Full name of fourth joint inventor, if any Shinichi AKAHANE	
第四共同発明者の署名	日付	Shinichi Akahane		Fourth inventor's signature	Date 10/27/00
住所	Residence Ebina, Japan JPX				
国籍	Citizenship Japan				
私書箱	Post Office Address c/o Hitachi, Ltd., Intellectual Property Group New Marunouchi Bldg. 5-1, Marunouchi 1-chome, Chiyoda-ku, Tokyo 100-8220, Japan				
第五共同発明者		5th		Full name of fifth joint inventor, if any Noboru TANABE	
第五共同発明者の署名	日付	Noboru Tanabe		Fifth inventor's signature	Date 10/23/00
住所	Residence Hadano, Japan JPX				
国籍	Citizenship Japan				
私書箱	Post Office Address c/o Hitachi, Ltd., Intellectual Property Group New Marunouchi Bldg. 5-1, Marunouchi 1-chome, Chiyoda-ku, Tokyo 100-8220, Japan				

(第六以降の共同発明者についても同様に記載し、署名をすること)

(Supply similar information and signature for sixth and subsequent joint inventors.)

第六共同発明者		Full name of sixth joint inventor, if any	
10/19/00 Yoshihito SAKO		Yoshihito SAKO	
第六共同発明者の署名	日付	Sixth inventor's signature	Date
		Yoshihito Sako	10/19/00
住所		Residence	
		Hadano, Japan JPX	
国籍		Citizenship	
		Japan	
私書箱		Post Office Address	
		c/o Hitachi, Ltd., Intellectual Property Group New Marunouchi Bldg. 5-1, Marunouchi 1-chome, Chiyoda-ku, Tokyo 100-8220, Japan	
第七共同発明者		Full name of seventh joint inventor, if any	
第七共同発明者の署名	日付	Seventh inventor's signature	Date
住所		Residence	
国籍		Citizenship	
私書箱		Post Office Address	
第八共同発明者		Full name of eighth joint inventor, if any	
第八共同発明者の署名	日付	Eighth inventor's signature	Date
住所		Residence	
国籍		Citizenship	
私書箱		Post Office Address	
第九共同発明者		Full name of ninth joint inventor, if any	
第九共同発明者の署名	日付	Ninth inventor's signature	Date
住所		Residence	
国籍		Citizenship	
私書箱		Post Office Address	

・ (第十以降の共同発明者についても同様に記載し、署名をすること)

(Supply similar information and signature for tenth and subsequent joint inventors.)